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## TUT-01: Age of Information as a New Data Freshness Metric in the IoT Era: From Theory to Implementation

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香港中文大學

# Outline of the Tutorial

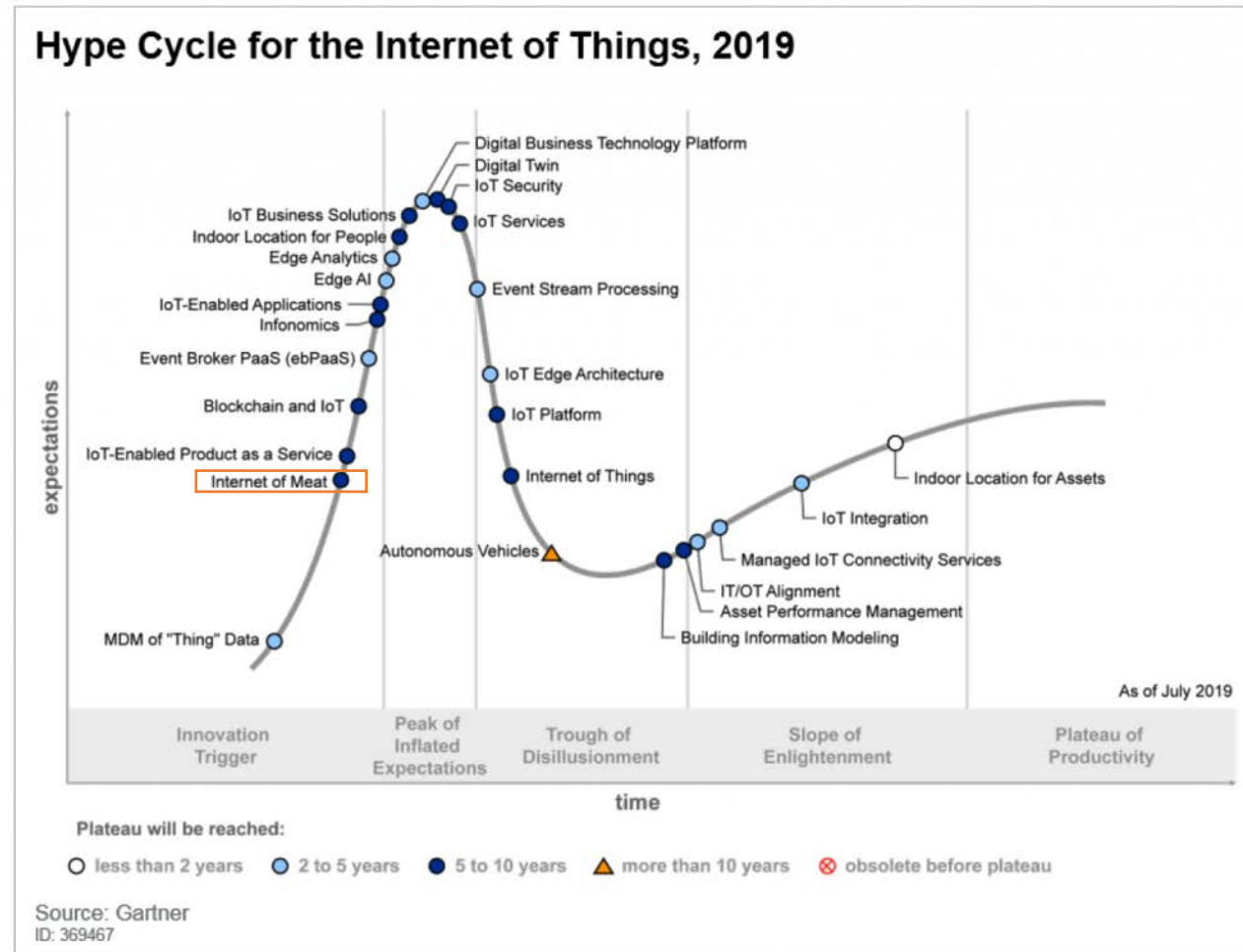
- Overview
- Modeling and Optimization of Aol
- Aol-oriented Multiuser Scheduling
- Aol-oriented Random Access
- Prototyping Testbed for Validation and Evaluation of Aol-oriented Designs
- Concluding remarks

# Internet of Things (IoT)



“**Wolf is coming**”  
for too many times

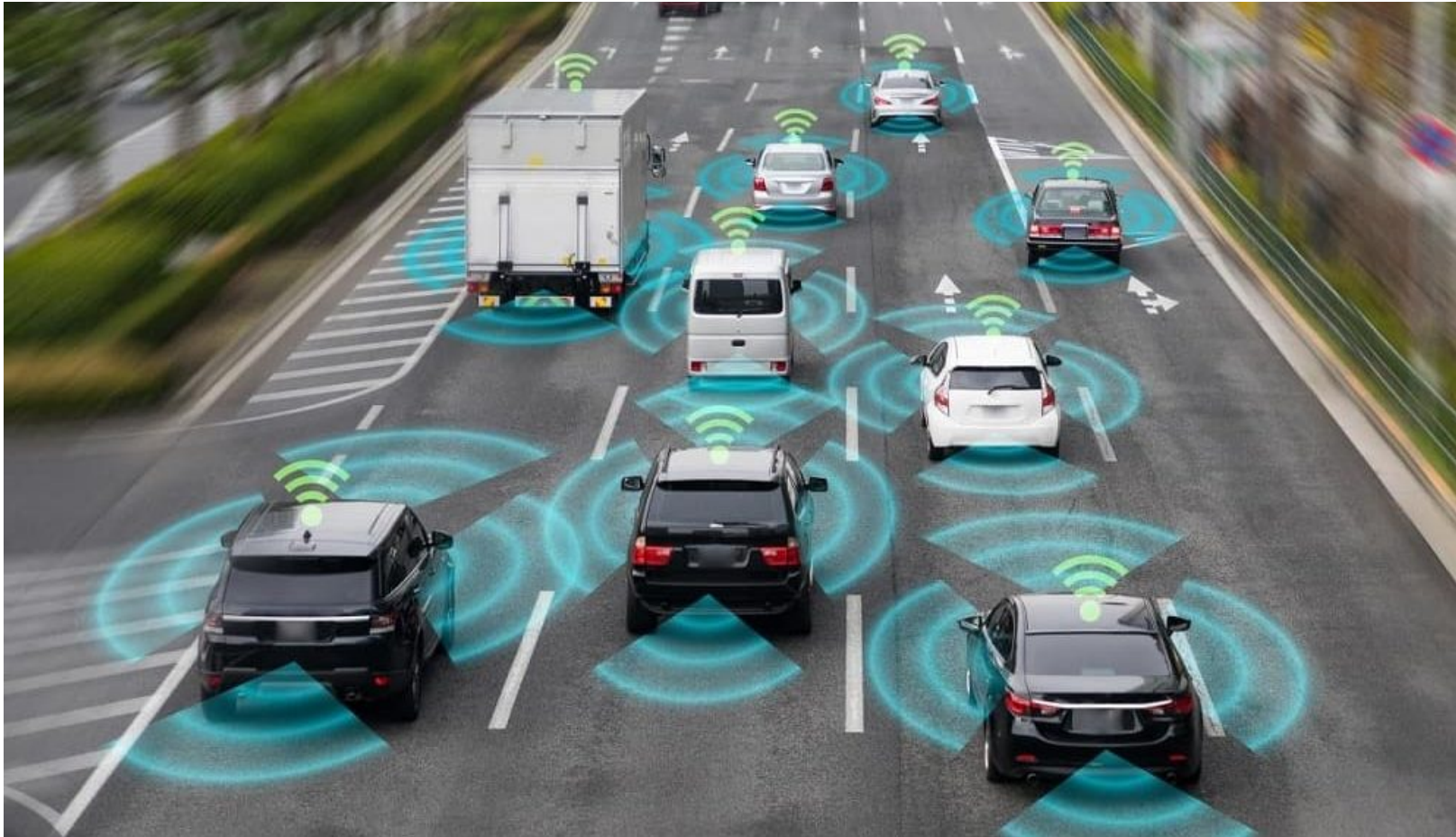
# IoT is really coming



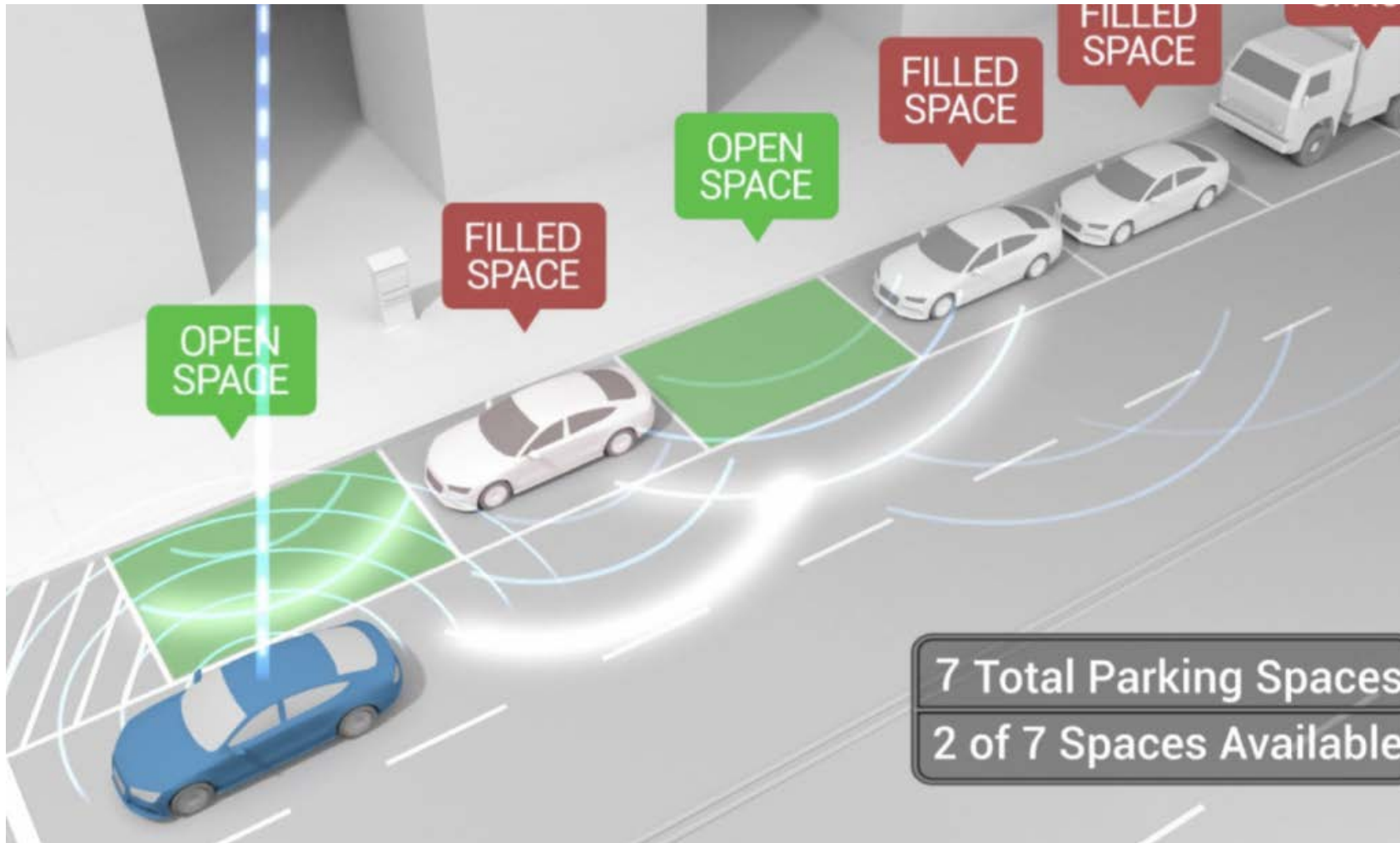
Source: <https://www.gartner.com/smarterwithgartner/3-major-trends-drive-gartner-hype-cycle-midsize-enterprises-2019/>



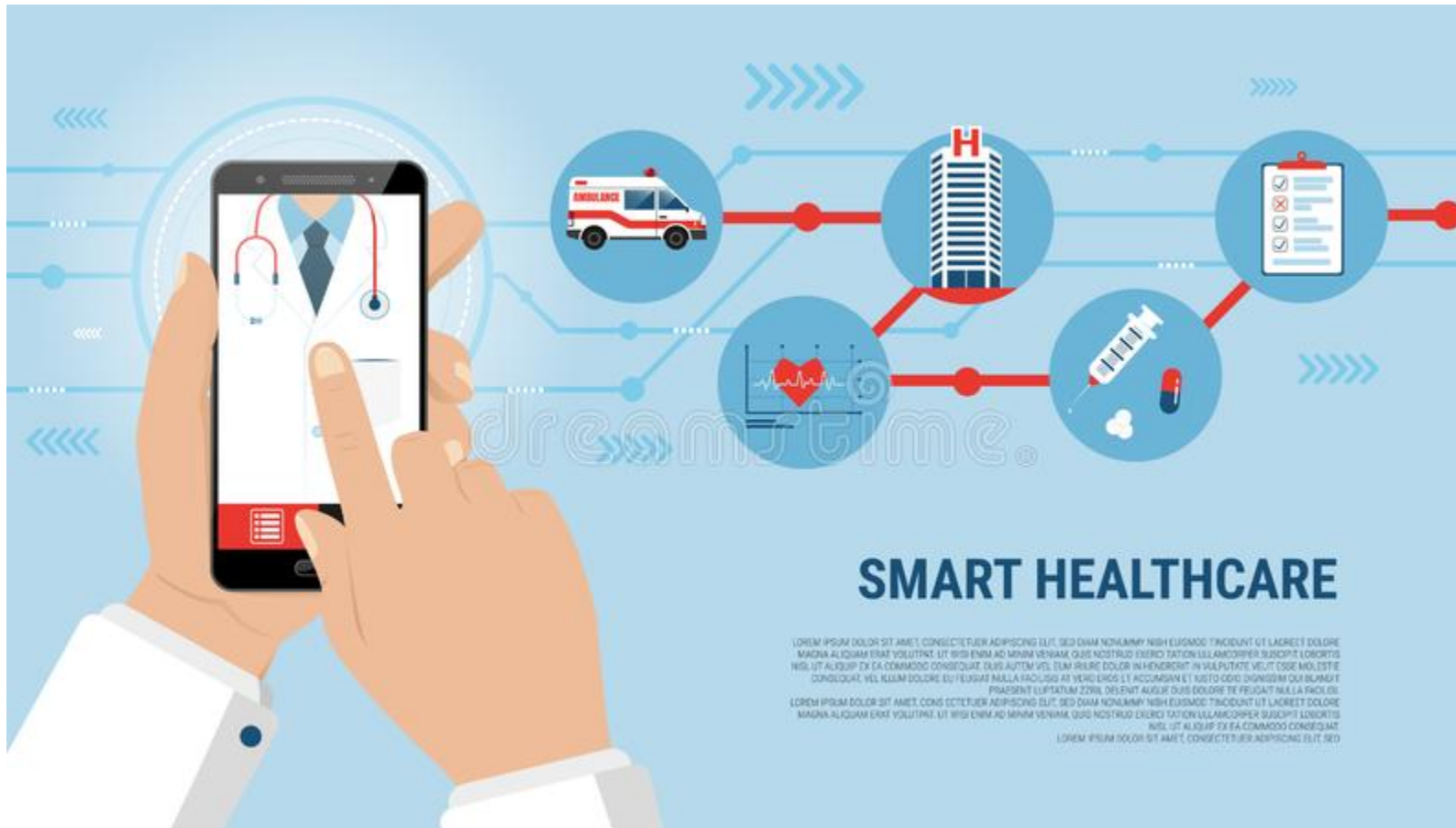
# Smart Transportation



# Smart Parking (Part of Smart City)



# Smart Healthcare





# Smart Factory



# Remote Monitoring

- All aforementioned IoT applications involve **remote monitoring**
- Various sensors monitor a certain **process** and send its sampled **statuses** of the monitored to a remote monitor --- status update system
- **Information** usually has the **highest value** when it is **fresh**.
- How to **quantify information freshness**?

# Age of Information as a New Data Freshness Metric in the IoT Era: From Theory to Implementation

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# Why we need fresh data

- Performance metrics used in the literature to characterize time sensitive information:
  - **Packet delay** tracks the time that elapsed from the generation of the packet until its delivery,
  - **inter-delivery time** is the time between two successive deliveries.
- These metrics are not sufficient to maintain fresh information at the destination.

# Why we need fresh data

- **Timeliness of status updates has emerged as a new field of network research.**
- **Even in the simplest queueing systems, timely updating is not the same as maximizing the utilization of the system that delivers these updates, nor the same as ensuring that updates are received with minimum delay.**
  - **While utilization is maximized by sending updates as fast as possible, this strategy will lead to a monitor receiving delayed updates that were backlogged in the communication system.**
  - **In this case, the timeliness of status updates at the receiver can be improved by reducing the update rate.**
  - **On the other hand, reducing the update rate will cause outdated status information at the receiver due to the lack of updates.**

S. Kaul, R. Yates, and M. Gruteser, "Real-time status: How often should one update?" IEEE INFOCOM 2012



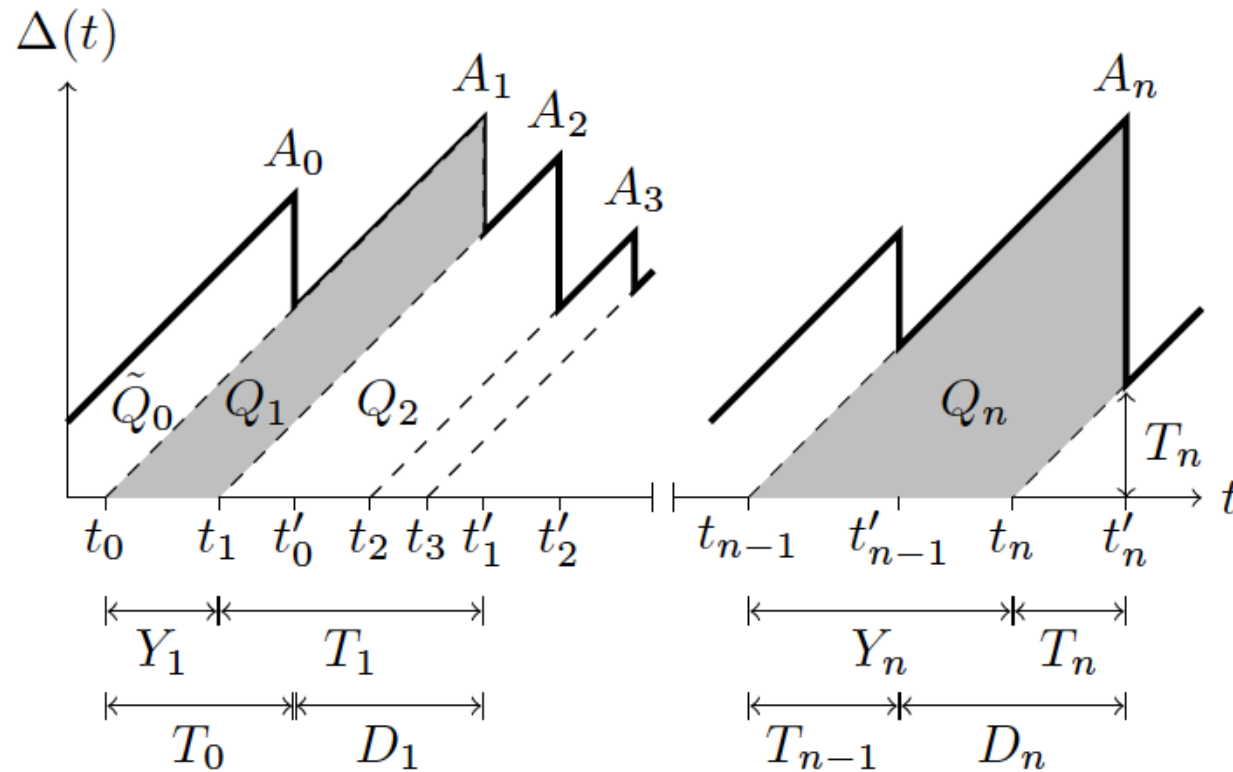
# Modeling and Optimization of Aol

Nikolaos Pappas

# Definition of Age of Information (AoI)

- AoI is an end-to-end metric that can be used to characterize latency in status updating systems and applications and captures the timeliness of the information.
- An update packet with timestamp  $u$  is said to have age  $t-u$  at a time  $t \geq u$ .
- An update is said to be fresh when its timestamp is the current time  $t$  and its age is zero.
- When the monitor's freshest received update at time  $t$  has timestamp  $u(t)$ , the age is the random process  $\Delta(t) = t - u(t)$ .

# Time Average Aol – Sawtooth Sample path



- $t_0, t_1, t_2, \dots$  times that are updates are generated
- $t'_0, t'_1, t'_2, \dots$  times that updates are received at the monitor
- For the  $n$ -th received update
  - $Y_n = t_n - t_{n-1}$  interarrival time
  - $T_n$  system time
  - $D_n = t'_n - t_{n-1}$  interdeparture time
  - $A_n$  corresponding peak age

A. Kosta, N. Pappas, V. Angelakis, "Age of Information: A New Concept, Metric, and Tool", Foundations and Trends in Networking: Vol. 12, No. 3, 2017.

R. D. Yates, Y. Sun, D. R. Brown III, S. K. Kaul, E. Modiano, and S. Ulukus, "Age of Information: An Introduction and Survey," arXiv:2007.08564, Jul. 2020

# Time Average Aol

$$\frac{1}{\mathcal{T}} \int_0^{\mathcal{T}} \Delta(t) dt$$

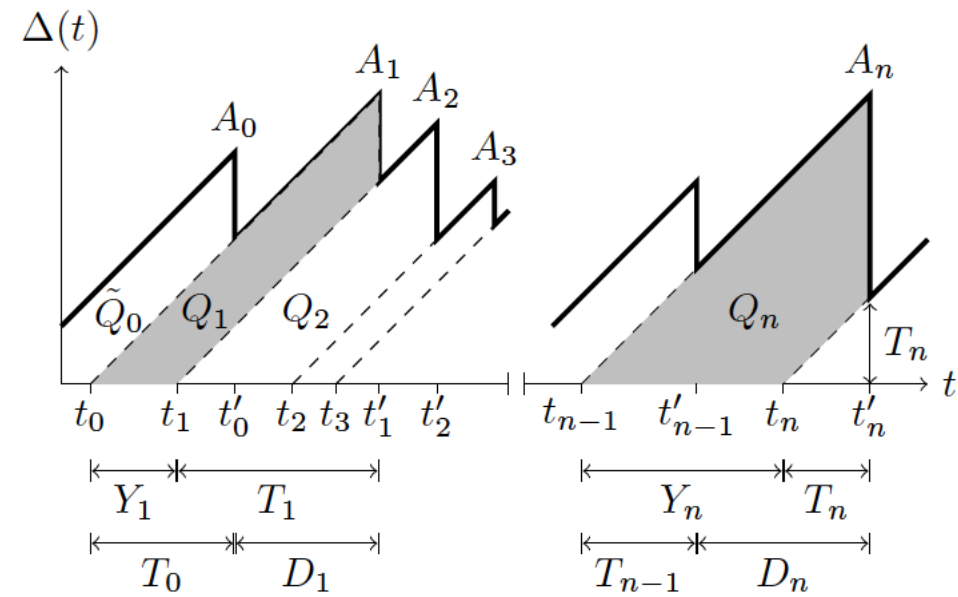
$$\mathcal{T} = t'_n$$

$$Q_n = \frac{1}{2}(T_n + Y_n)^2 - \frac{1}{2}T_n^2 = Y_n T_n + Y_n^2/2$$

$$N(\mathcal{T})/\mathcal{T} \rightarrow 1/E[Y]$$

$$\frac{1}{N(\mathcal{T})} \sum_{j=1}^{N(\mathcal{T})} Q_j \rightarrow E[Q] \quad \mathcal{T} \rightarrow \infty$$

$$\Delta = \frac{E[Q_n]}{E[Y_n]} = \frac{E[Y_n T_n] + E[Y_n^2]/2}{E[Y_n]}$$



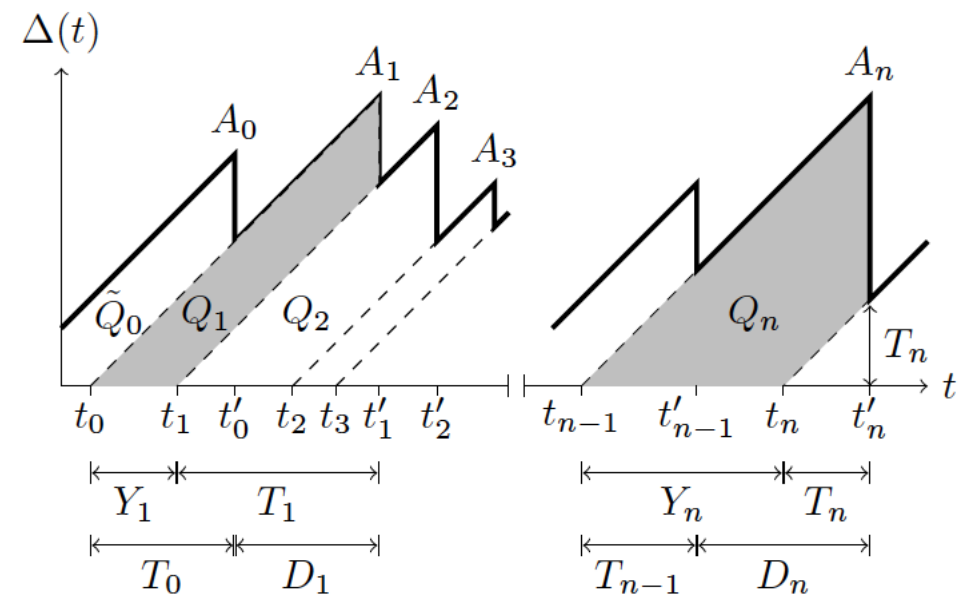
**Time Average Aol**  $\frac{1}{\mathcal{T}} \int_0^{\mathcal{T}} \Delta(t) dt$   $\mathcal{T} = t'_n$

$$Q_n = \frac{1}{2}(T_n + Y_n)^2 - \frac{1}{2}T_n^2 = Y_n T_n + Y_n^2/2$$

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$$\Delta = \frac{E[Q_n]}{E[Y_n]} = \frac{E[Y_n T_n] + E[Y_n^2]/2}{E[Y_n]}$$



Large interarrival time allows queue to be empty, thus, the waiting time can be small, causing small system time  $T_n$ .  $Y_n$  and  $T_n$  are negatively correlated which complicates the calculation of  $E[Y_n T_n]$

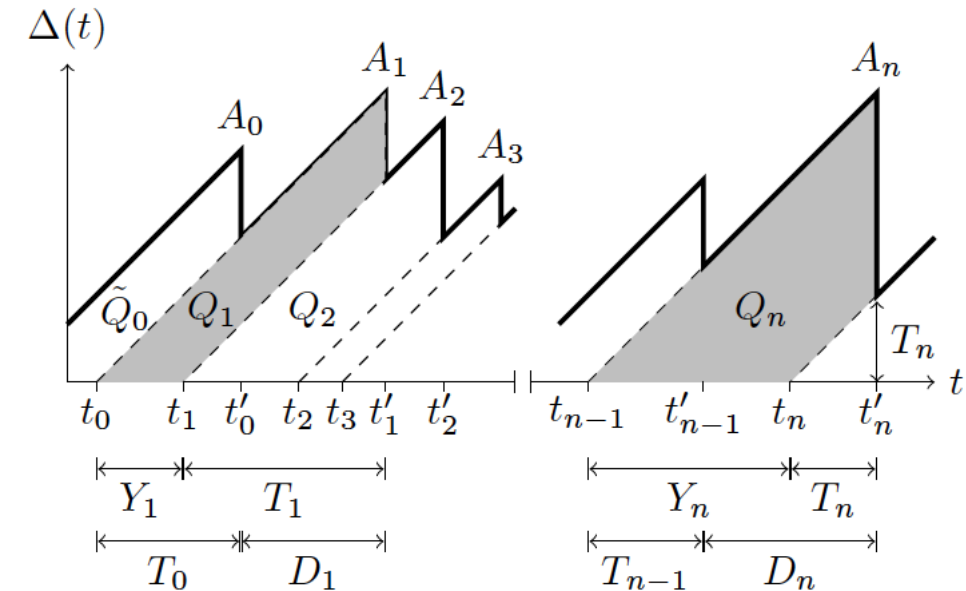
# Peak AoI

- Alternative and more tractable metric than AoI

$$\Delta^{(p)} = \lim_{\mathcal{T} \rightarrow \infty} \frac{1}{N(\mathcal{T})} \sum_{n=1}^{N(\mathcal{T})} A_n$$

$$\Delta^{(p)} = \mathbf{E}[A] = \mathbf{E}[T_{n-1}] + \mathbf{E}[D_n]$$

$$\Delta^{(p)} = \mathbf{E}[Y] + \mathbf{E}[T]$$



M. Costa, M. Codreanu, and A. Ephremides, "Age of information with packet management," IEEE ISIT 2014.

# Non-linear Ageing

- AoI grows over time linearly, the performance degradation caused by information aging may not be a linear function of time.
- One way to capture the nonlinear behavior of information aging is to define freshness and staleness as nonlinear functions of AoI.
- A penalty function of the AoI is non- decreasing. This requirement is aligned with the fact that outdated data is usually less desirable than fresh data.

Y. Sun, E. Uysal-Biyikoglu, R. Yates, C. E. Koksall, and N. B. Shroff, "Update or wait: How to keep your data fresh," IEEE Trans. Inf. Theory, 2017.

A. Kosta, N. Pappas, A. Ephremides, and V. Angelakis, "Age and value of information: Non-linear age case," IEEE ISIT 2017.

Y. Sun and B. Cyr, "Sampling for data freshness optimization: Nonlinear age functions," JCN 2019.

A. Kosta, N. Pappas, A. Ephremides, and V. Angelakis, "The cost of delay in status updates and their value: Non-linear ageing," Accepted, IEEE TCOM 2020.

# Cost of Update Delay (CoUD)

- CoUD metric associates the cost of staleness with the statistics of the source
- $C(t) = f_s(t-u(t))$ 
  - $f_s(t)$  is a monotonically increasing function
  - $u(t)$  timestamp of the most recently received update
- Different cost functions can represent different utilities

A. Kosta, N. Pappas, A. Ephremides, V. Angelakis, "Age and Value of Information: Non-linear Age Case", *IEEE ISIT* 2017.

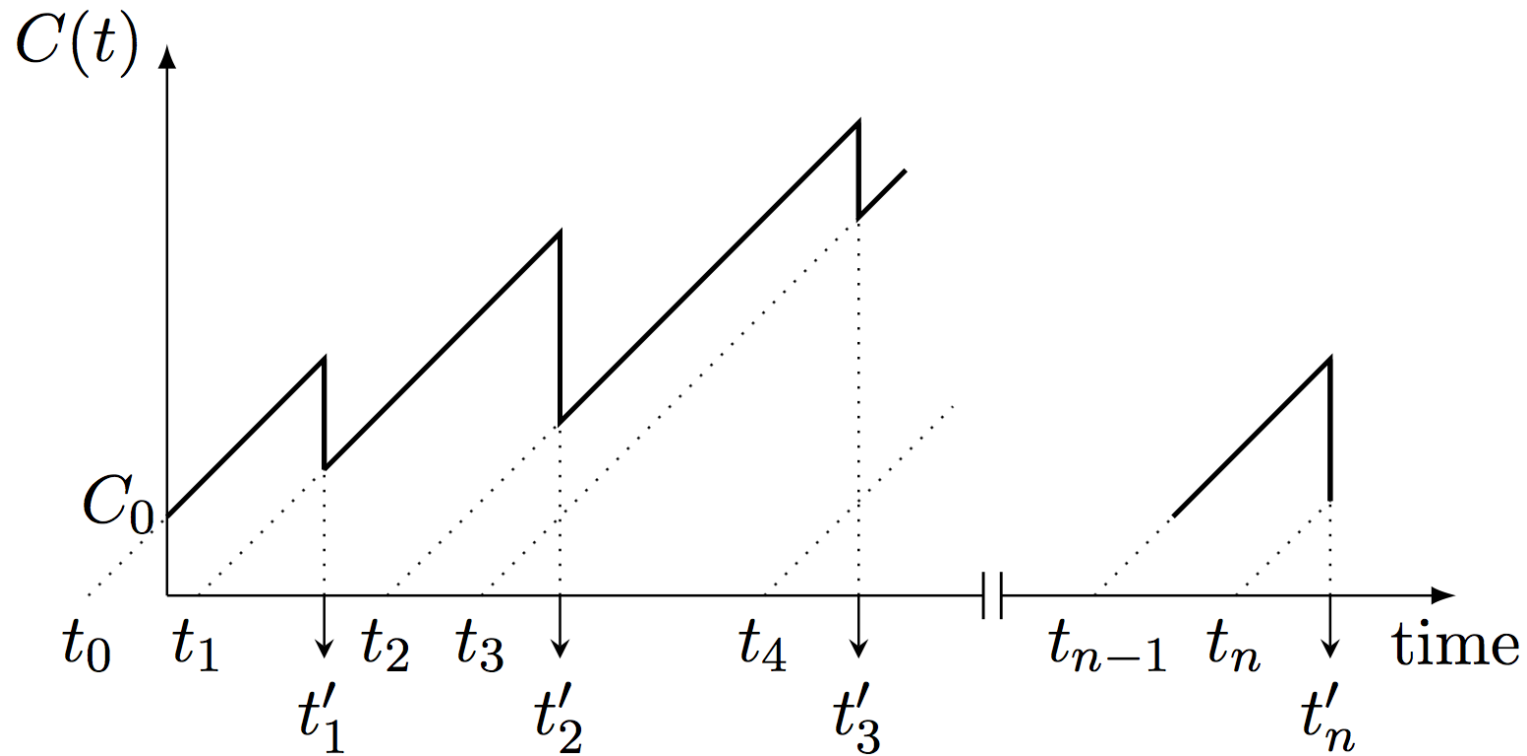
A. Kosta, N. Pappas, A. Ephremides, V. Angelakis, "The Cost of Delay in Status Updates and their Value: Non-linear Ageing", *Accepted*, *IEEE TCOM*, 2020.



# Cost of Update Delay (CoUD): The linear case

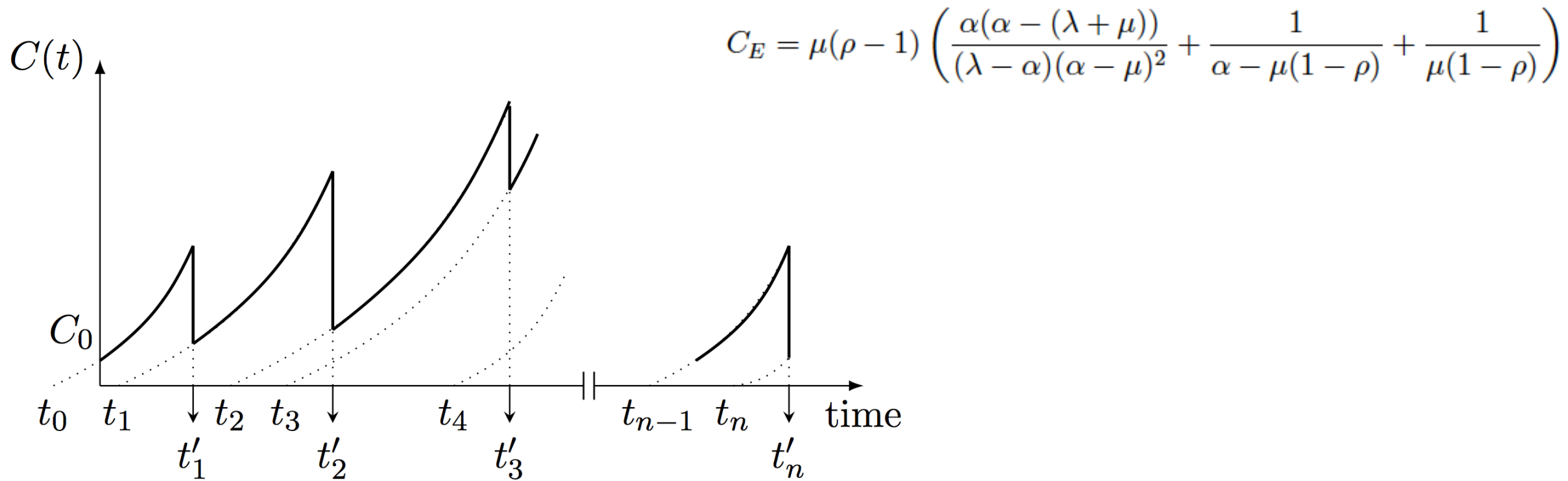
$$f_s(t) = \alpha t$$

$$C_P = \alpha \frac{1}{\mu} \left( 1 + \frac{1}{\rho} + \frac{\rho^2}{1 - \rho} \right)$$



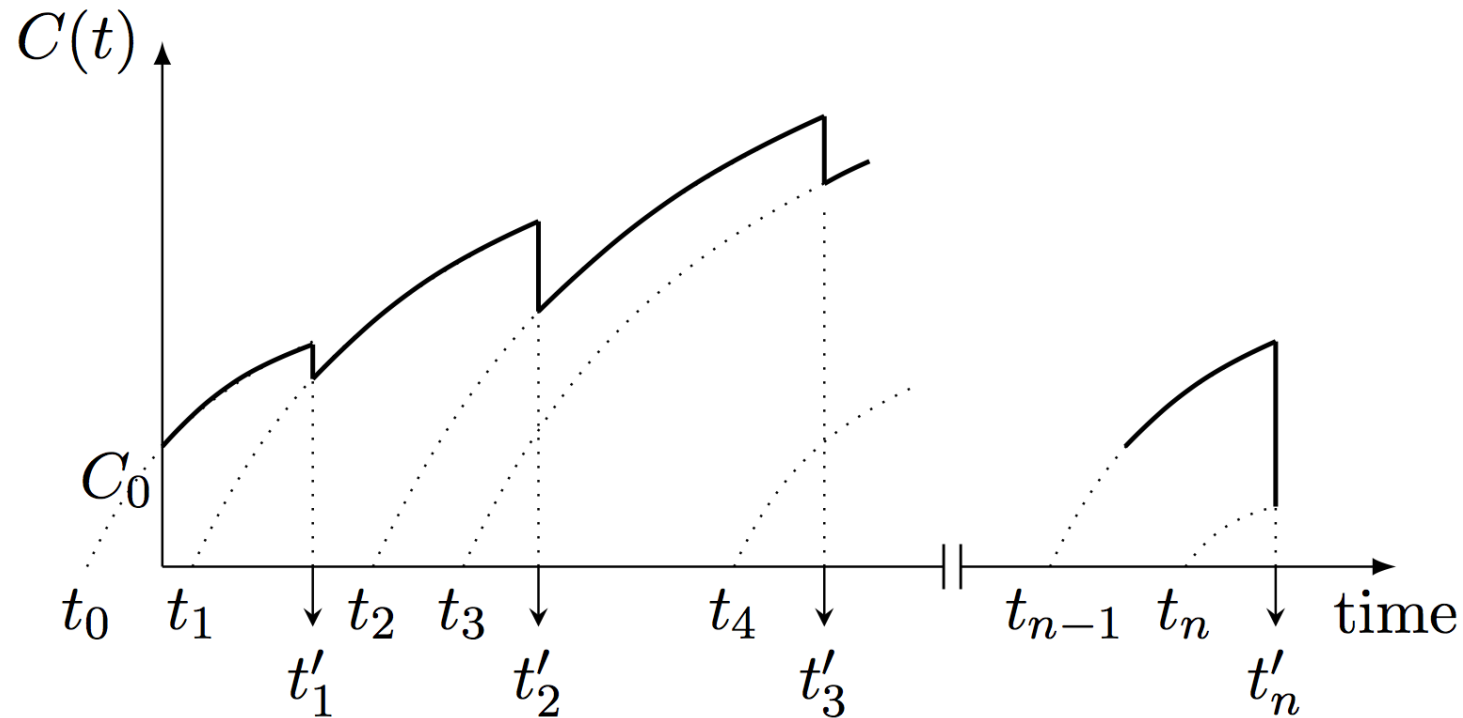
# Cost of Update Delay (CoUD): The exponential case

$$f_s(t) = e^{\alpha t} - 1 \longleftrightarrow \text{low autocorrelation}$$



# Cost of Update Delay (CoUD): The logarithmic case

$f_s(t) = \log(\alpha t + 1) \longleftrightarrow$  high autocorrelation



$$C_L = \frac{1}{\alpha(\lambda - \mu)^2} \left( e^{-\frac{\mu\rho}{\alpha}} \left( \mu(1 - \rho) Ei \left[ -\frac{\mu}{\alpha} \right] (\alpha\mu + \lambda^2 - \lambda\mu) e^{\frac{\mu(\rho+1)}{\alpha}} - \alpha\mu^2(1 - \rho) Ei \left[ -\frac{\lambda}{\alpha} \right] e^{\frac{\lambda+\mu\rho}{\alpha}} - \alpha e^{\mu/\alpha} (\lambda - \mu)^2 Ei \left[ -\frac{\mu(1 - \rho)}{\alpha} \right] \right) - \alpha\lambda(1 - \rho)(\mu - \lambda) \right)$$

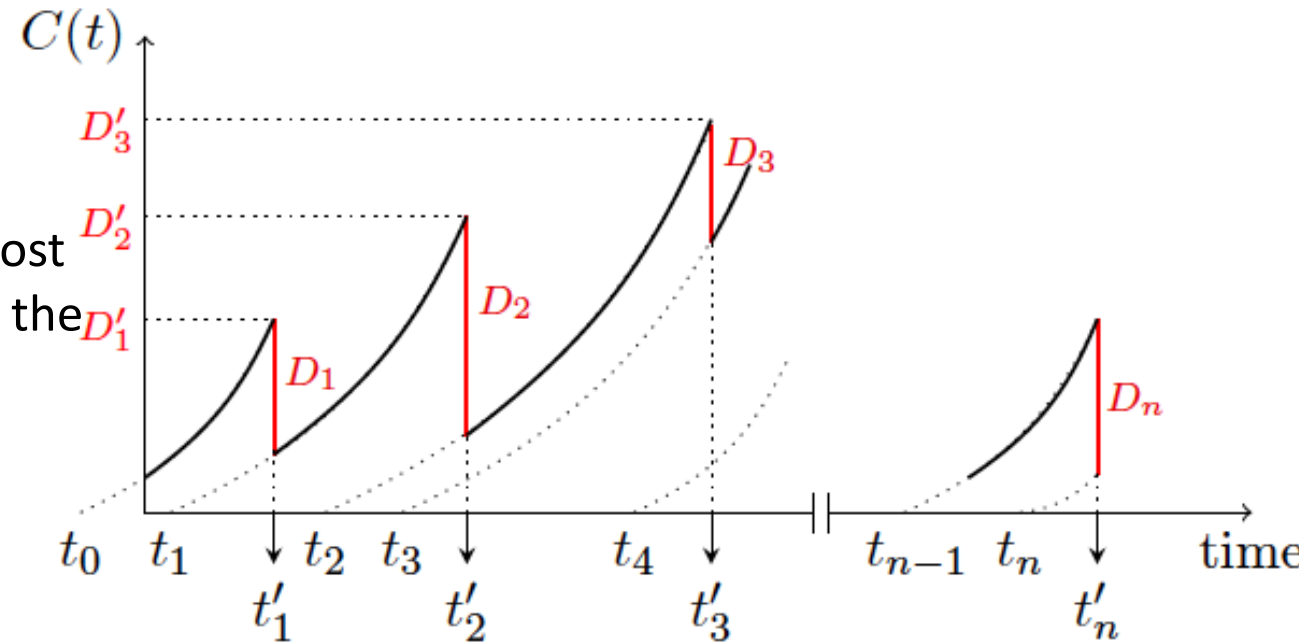
# Value of Information of update (VoIU)

- It captures the degree of importance of an update

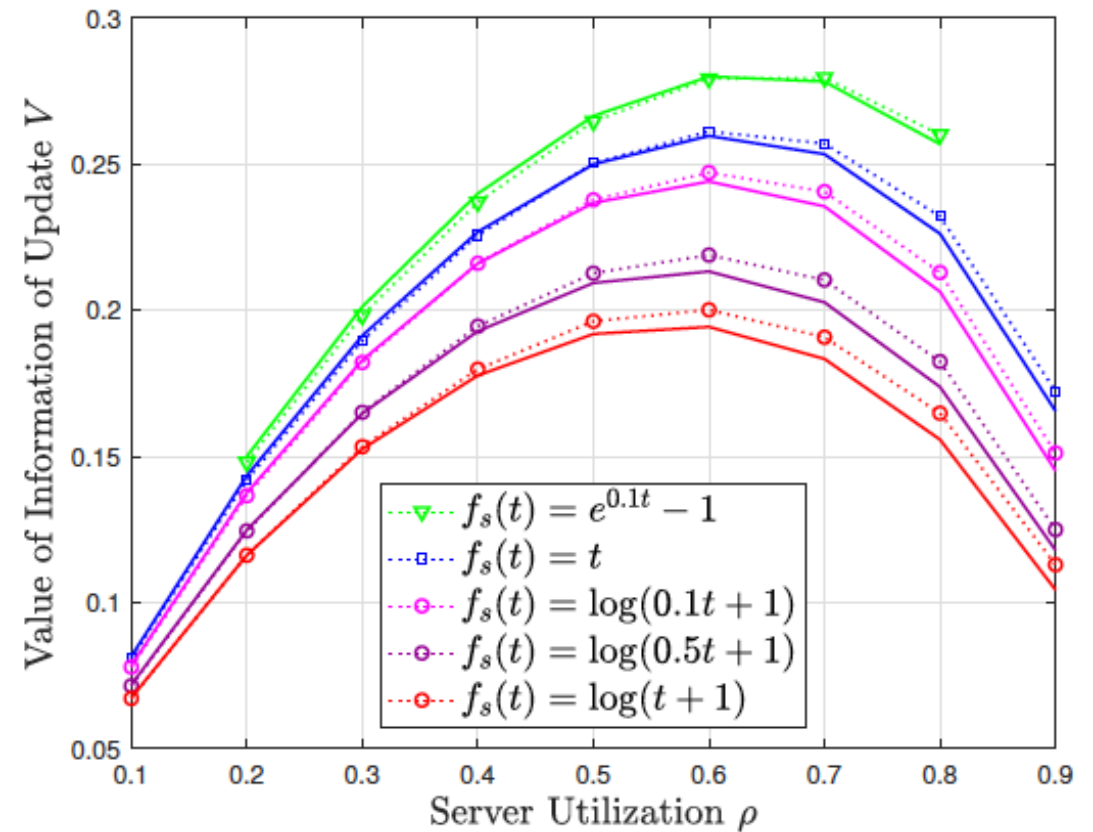
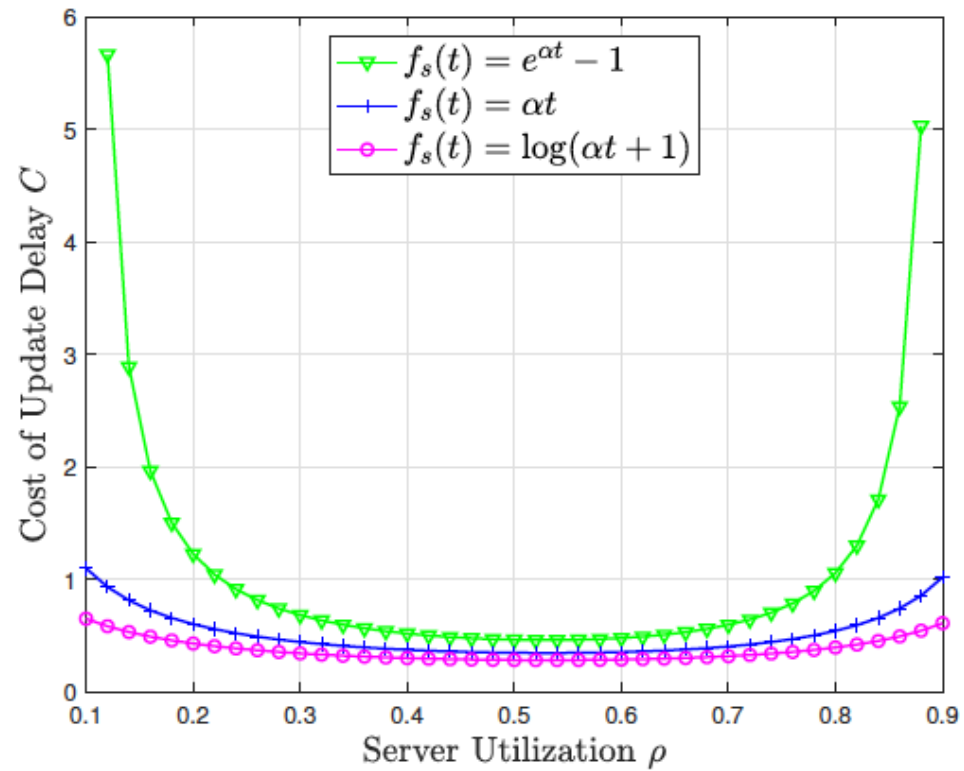
$$V_i = \frac{f_s(t'_i - t_{i-1}) - f_s(t'_i - t_i)}{f_s(t'_i - t_{i-1})} = \frac{D_i}{D'_i}$$

In the linear CoUD case, VoIU is independent of the cost assigned per time unit → the Value is independent of the slope.

$$V_i = \lim_{t'_i \rightarrow t_i} \frac{f_s(t'_i - t_{i-1}) - f_s(t'_i - t_i)}{f_s(t'_i - t_{i-1})} = 1$$



# Numerical evaluation



# Single-source and single-server systems

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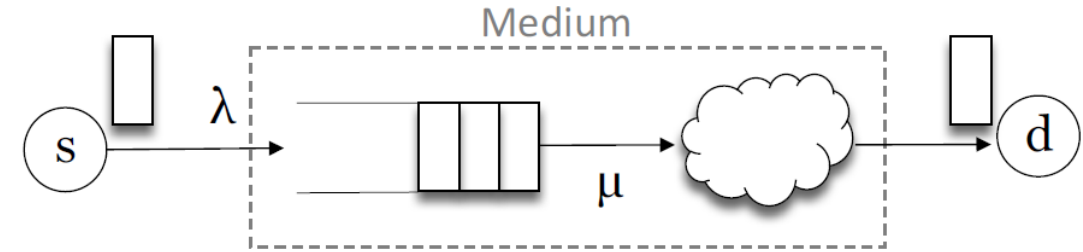
*i.i.d interarrival times with expected value  $E[Y]$*

$\lambda=1/ E[Y]$  : arrival rate

$E[S]$  : expected service time

$\mu=1/ E[S]$  : service rate

$\rho = \lambda/ \mu$  : offered load



For FCFS M/M/1 queue the average is  $\Delta_{M/M/1} = \frac{1}{\mu} \left( 1 + \frac{1}{\rho} + \frac{\rho^2}{1 - \rho} \right)$

The optimal age is achieved for  $\rho^* \approx 0.53$

- Optimal age is achieved by choosing a  $\lambda$  which makes the server being slightly busy than being idle.
- If  $\rho$  is close to 1 we maximize the throughput, which is the number of packets delivered to the destination per time.
- If  $\rho$  is close to 0, we minimize the delay.

S. Kaul, R. Yates, and M. Gruteser, "Real-time status: How often should one update?" IEEE INFOCOM 2012.

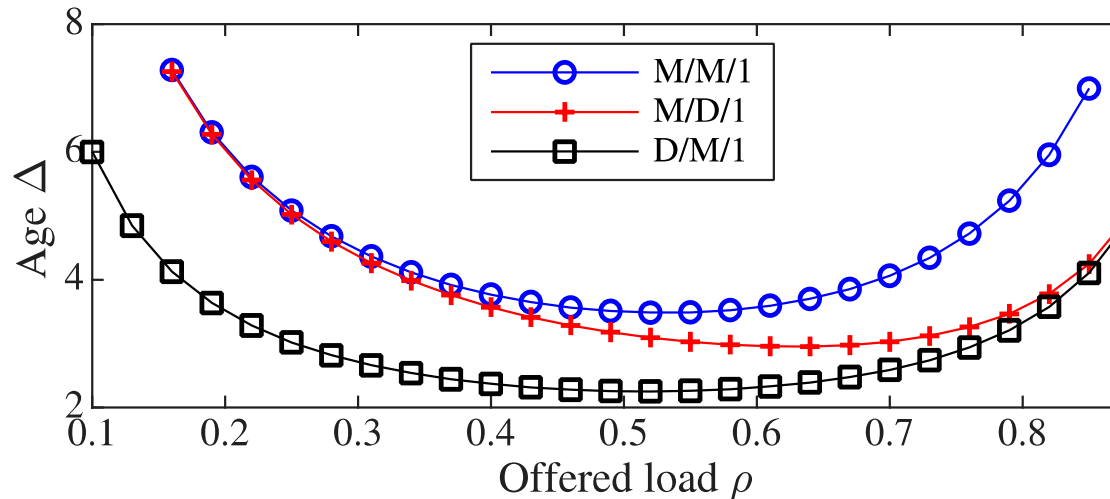
M. Costa, M. Codreanu, and A. Ephremides, "On the age of information in status update systems with packet management," IEEE Trans. Info. Theory 2016.

Y. Inoue, H. Masuyama, T. Takine, and T. Tanaka, "A general formula for the stationary distribution of the age of information and its application to single-server queues," IEEE Trans. Info. Theory 2019.

# Single-source and single-server systems

For M/D/1 and D/M/1 queues the average AoI are given by

$$\Delta_{M/D/1} = \frac{1}{\mu} \left( \frac{1}{2(1-\rho)} + \frac{1}{2} + \frac{(1-\rho)\exp(\rho)}{\rho} \right) \quad \Delta_{D/M/1} = \frac{1}{\mu} \left( \frac{1}{2\rho} + \frac{1}{1-\gamma(\rho)} \right) \quad \gamma(\rho) = -\rho \mathcal{W} \left( -\rho^{-1} e^{(-1/\rho)} \right)$$



- At low load, randomness in the interarrivals dominates the average age.
- At high load, M/D/1 and D/M/1 substantially outperform M/M/1 because the determinism in either arrivals or service helps to reduce the average queue length.
- For each queue, we observe a unique value of  $\rho$  that minimizes the average age.

# Single-source and single-server systems – Packet management

- The arrival rate can be optimized to balance frequency of updates against congestion.
- Study of lossy queues that may discard an arriving update while the server was busy or replace an older waiting update with a fresher arrival.
- Packet management:
  - M/M/1/1 queue that blocks and clears a new arrival while the server is busy,
  - M/M/1/2 queue that will queue one waiting packet but blocks an arrival when the waiting space is occupied,
  - M/M/1/2\* queue

S. Kaul, R. Yates, and M. Gruteser, "Status updates through queues," CISS 2012.

N. Pappas, J. Gunnarsson, L. Kratz, M. Kountouris, V. Angelakis, "Age of Information of Multiple Sources with Queue Management", IEEE ICC 2015.

M. Costa, M. Codreanu, and A. Ephremides, "On the age of information in status update systems with packet management," IEEE Trans. Info. Theory 2016.

A. Kosta, N. Pappas, A. Ephremides, V. Angelakis, "Age of Information Performance of Multiaccess Strategies with Packet Management", IEEE/KICS JCN, June 2019.



# Multiple sources at a single server source

Each source  $i$  generates updates as an independent Poisson process of rate  $\lambda_i$  and the service time  $S$  of an update has expected value  $1/\mu$ . The total offered load is  $\rho = \sum_i \lambda_i / \mu$ .

R. D. Yates and S. K. Kaul, "The age of information: Real-time status updating by multiple sources," IEEE Trans. Info. Theory, 2019.

L. Huang and E. Modiano, "Optimizing age-of-information in a multiclass queueing system," IEEE ISIT 2015.

E. Najm and E. Telatar, "Status updates in a multi-stream M/G/1/1 preemptive queue," IEEE INFOCOM Workshops 2018.

M. Moltafet, M. Leinonen, and M. Codreanu, "On the age of information in multi-source queueing models," IEEE TCOM 2020.

$P_\lambda = \mathbb{E}[e^{-\lambda S}]$  Laplace transform of service time  $S$

$$\lambda = \lambda_1 + \dots + \lambda_N$$

For the user  $i$  in an M/G/1/1 system the average AoI is  $\Delta_i = \frac{1}{\lambda_i P_\lambda}$   
and Peak AoI is

$$\Delta_i^{(p)} = \frac{1}{\lambda_i P_\lambda} + \frac{\mathbb{E}[S e^{-\lambda S}]}{P_\lambda}$$

# Multiple sources at a single server source

Each source  $i$  generates updates as an independent Poisson process of rate  $\lambda_i$  and the service time  $S$  of an update has expected value  $1/\mu$ . The total offered load is  $\rho = \sum_i \lambda_i / \mu$ .

R. D. Yates and S. K. Kaul, "The age of information: Real-time status updating by multiple sources," IEEE Trans. Info. Theory, 2019.

L. Huang and E. Modiano, "Optimizing age-of-information in a multiclass queueing system," IEEE ISIT 2015.

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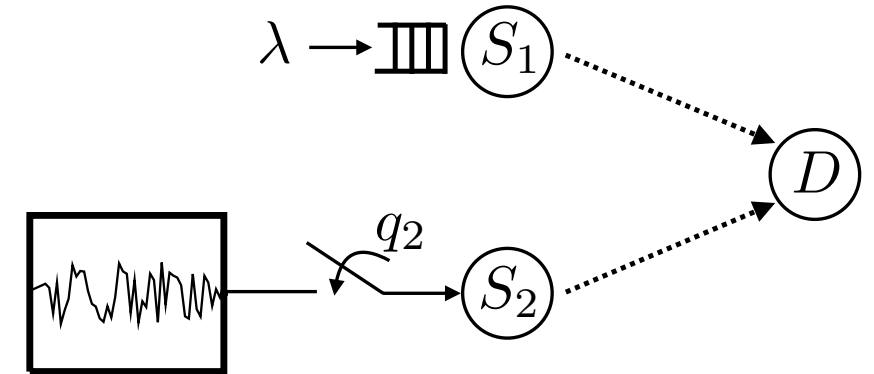
A case of heterogeneous traffic sharing the same queue will be discussed a bit later.

# Towards a complete characterization of the Aol distribution

- Stochastic hybrid systems are utilized to analyze Aol moments and the moment generating function of Aol in networks
  - R. D. Yates, “The Age of Information in Networks: Moments, Distributions, and Sampling,” *IEEE Trans. Info. Theory* 2020.
- A general formula of the stationary distribution of Aol is obtained and applied to a wide class of continuous-time single server queues with different disciplines
  - Y. Inoue, H. Masuyama, T. Takine, and T. Tanaka, “A general formula for the stationary distribution of the age of information and its application to single-server queues,” *IEEE Trans. Info. Theory* 2019.
- The distribution of Aol for the GI/GI/1/1 and GI/GI/1/2\* systems, under non-preemptive scheduling
  - J. P. Champati, H. Al-Zubaidy, and J. Gross, “On the distribution of aoi for the GI/GI/1/1 and GI/GI/1/2\* systems: Exact expressions and bounds,” *IEEE INFOCOM* 2019.
- The Aol distribution in bufferless systems
  - G. Kesidis, T. Konstantopoulos, and M. Zazanis, “The distribution of age-of-information performance measures for message processing systems,” *arXiv:1904.05924*, 2019.
- The complete characterization of the Aol stationary distribution in a discrete time queueing system for three cases: FCFS, preemptive LCFS, a bufferless system with packet dropping. *We provide a methodology for analyzing general non-linear age functions, using representations of functions as power series.*
  - A. Kosta, N. Pappas, A. Ephremides, V. Angelakis, "The Age of Information in a Discrete Time Queue: Stationary Distribution and Non-linear Age Mean Analysis", *arxiv*, June 2020. (shorter version in *IEEE ICC* 2020)

# AoI and Delay Violation Probability Interplay in the Two-user MAC

- Two sources sending packets to a common destination.
- Source  $S_1$  has external traffic with stringent delay requirements.
- Source  $S_2$  monitors a sensor and samples a status update on each slot w. p.  $q_2$ . Then, transmits the update to the destination through a channel with success probability  $p_2$ . If the transmission of a status update fails, then it is dropped.
- Time is slotted, fixed rate  $R$  transmissions, Rayleigh fading,
- Instantaneous and error-free ACK/NACK.
- $S_i$  transmits with power  $P_i$ ,  $i=1,2$ .



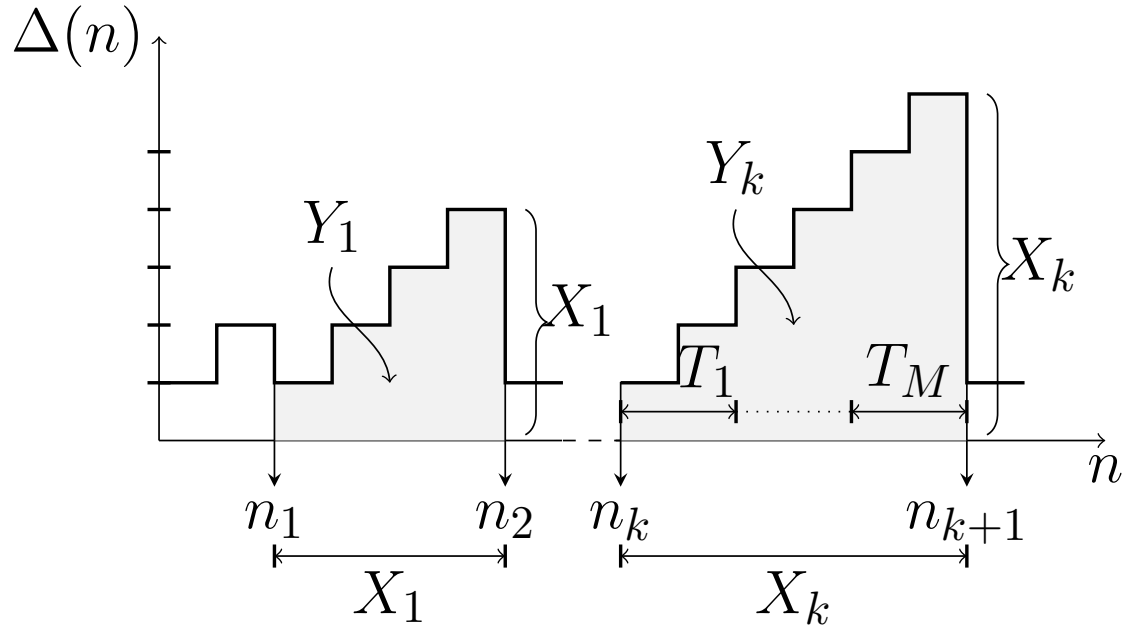
N. Pappas, M. Kountouris, "Delay Violation Probability and Age-of-information Interplay in the Two-user Multiple Access Channel", *IEEE SPAWC 2019*.

# Average AoI

$T_i$ : time between two consecutive attempted transmissions

$X_k$ : elapsed time at the destination between successful reception of  $k$ -th and the  $(k + 1)$ -th status updates

$M$ : number of attempted transmissions between two successfully received status updates at  $D$

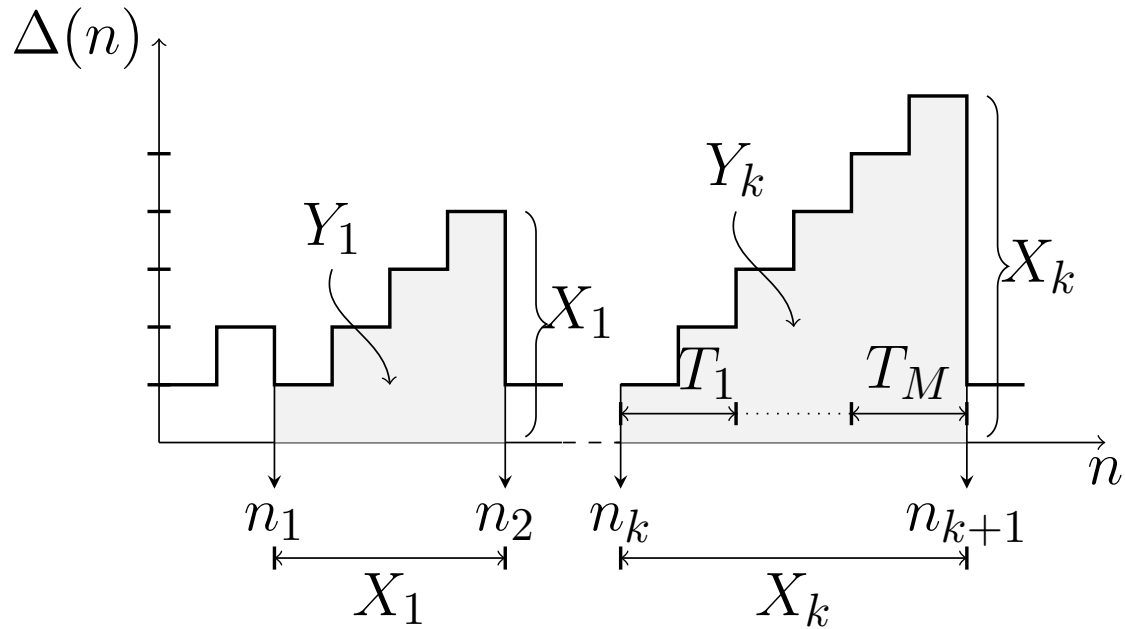


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$$X_k = \sum_{i=1}^M T_i \quad \Delta_N = \frac{1}{N} \sum_{n=1}^N \Delta(n) = \frac{1}{N} \sum_{k=1}^K Y_k = \frac{K}{N} \frac{1}{K} \sum_{k=1}^K Y_k$$

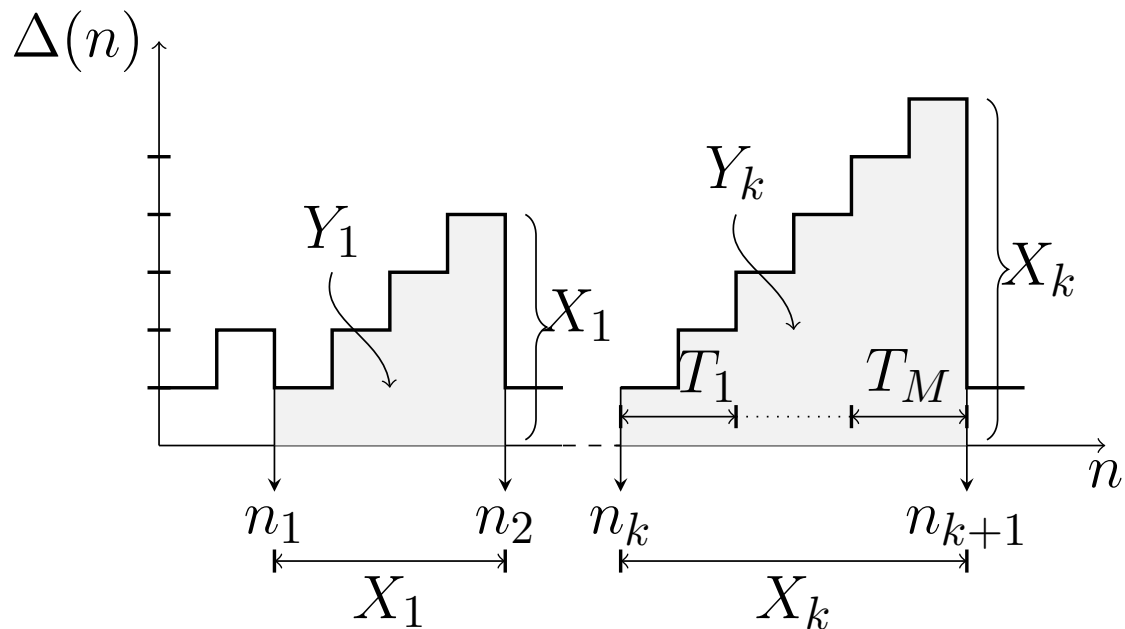
$$\Delta = \lim_{N \rightarrow \infty} \Delta_N = \frac{\mathbb{E}[Y]}{\mathbb{E}[X]} \quad Y_k = \sum_{m=1}^{X_k} m = \frac{X_k(X_k + 1)}{2}$$

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$$\Delta_N = \frac{K}{N} \frac{1}{K} \sum_{k=1}^K Y_k = \frac{\mathbb{E} \left[ \frac{X_k^2}{2} + \frac{X_k}{2} \right]}{\mathbb{E}[X]} = \frac{\mathbb{E}[X^2]}{2\mathbb{E}[X]} + \frac{1}{2}$$

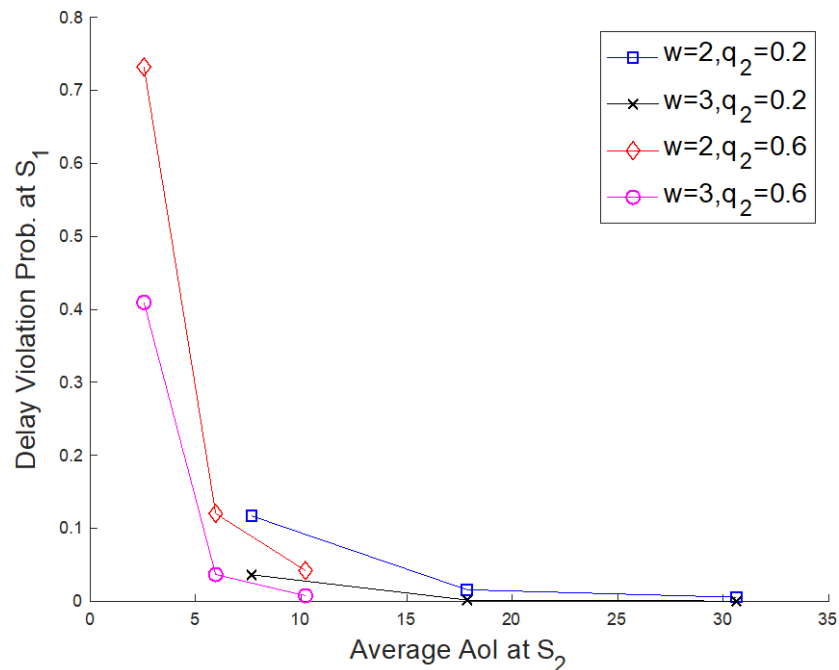
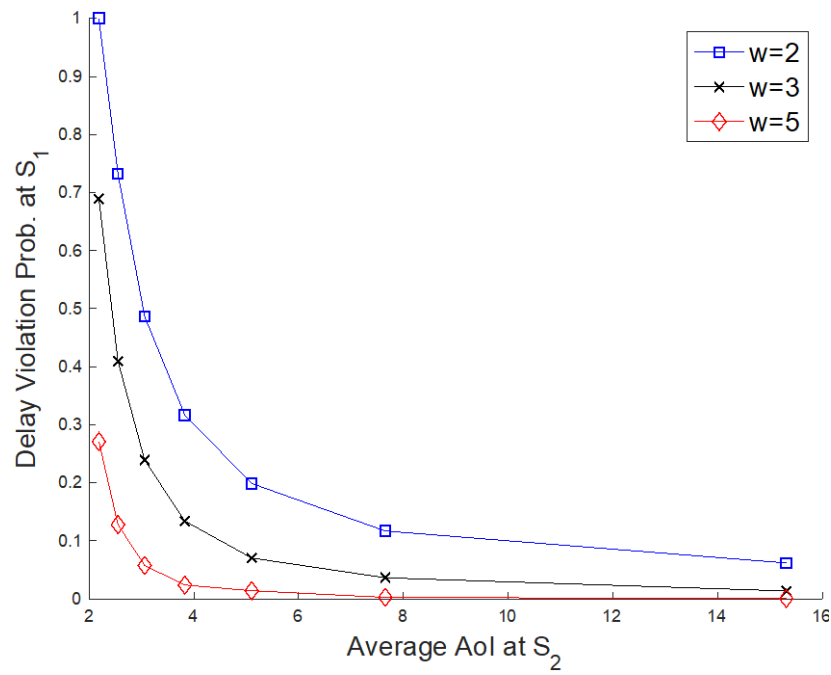
$$\mathbb{E}[X] = \sum_{M=1}^{\infty} M \mathbb{E}[T] (1 - p_2)^{M-1} p_2 = \frac{\mathbb{E}[T]}{p_2}$$

$$\mathbb{E}[X^2] = \sum_{M=1}^{\infty} \mathbb{E}[X^2 | M] (1 - p_2)^{M-1} p_2$$

$$\stackrel{p_2 > 0}{=} \frac{\mathbb{E}[T^2]}{p_2} + \frac{2(1 - p_2)\mathbb{E}[T]^2}{p_2^2}$$

$$\Delta = \frac{\mathbb{E}[T^2]}{2\mathbb{E}[T]} + \frac{\mathbb{E}[T](1 - p_2)}{p_2} + \frac{1}{2}$$

$$= \frac{1}{q_2 p_2}$$



- The latency constraint  $w$  does not affect the average Aol.
- As  $w$  increases, the delay violation probability decreases since  $S_1$  becomes more delay tolerant.
- Increasing the transmit prob. results in significant decrease of the delay violation probability and an increase of Aol due to larger interference.

### *Main takeaway*

*Both delay violation probability and Aol can be kept low even for stringent delay constraints if the sampling rate is properly adapted.*



## Zero-wait

Y. Sun, E. Uysal-Biyikoglu, R. Yates, C. E. Koksal, and N. B. Shroff, “Update or wait: How to keep your data fresh,” *IEEE INFOCOM, 2016* and *IEEE Trans. Inf. Theory, 2017*.

## Queueing Networks

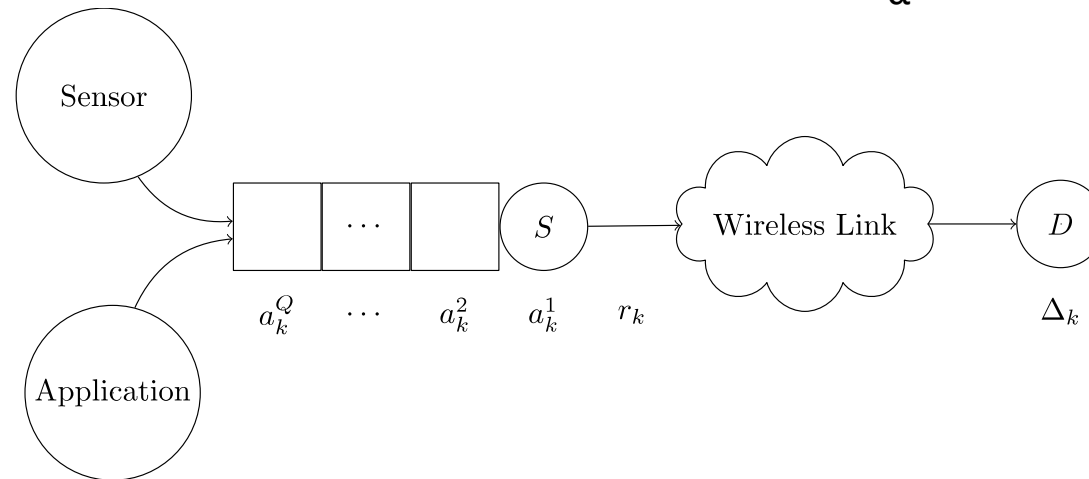
C. Kam, S. Kompella, and A. Ephremides, “Age of information under random updates,” *IEEE ISIT 2013*.

C. Kam, S. Kompella, G. D. Nguyen, and A. Ephremides, “Effect of message transmission path diversity on status age,” *IEEE Trans. Info. Theory, 2016*.

R. D. Yates, “Status updates through networks of parallel servers,” *IEEE ISIT 2018*.

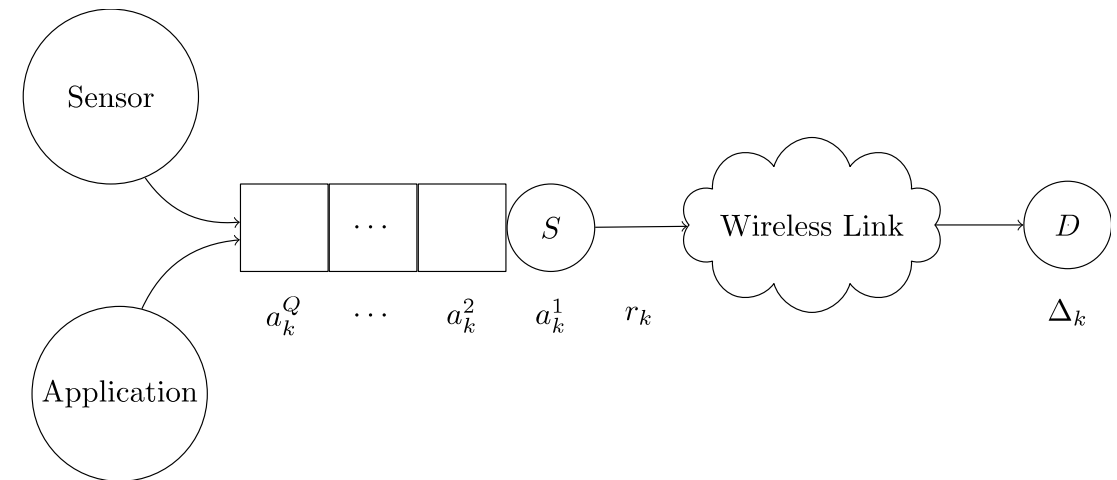
# Optimal status update generation in a device with heterogeneous traffic

- A source node transmitting data to a destination through a wireless link.
- Source node consists of a sensor and an application and has a finite capacity queue
  - Sensor is generating status updates in forms of packets
  - Application generates data packets
- Discrete time
- Wireless link with success probability  $P_s$  and the number of retransmissions is fixed, equal to  $r_{\max}$
- Application generates a data packet in each slot with prob.  $P_\alpha$



# Optimal status update generation in a device with heterogeneous traffic

- The source decides in each slot if a status update from the sensor is generated or not
- All packets generated in a slot are enqueued unless the queue is full
- The source must satisfy a hard constraint  $\Delta_{\max}$  on age of information for the sensor
- When the age reaches the value  $\Delta_{\max}$  then
  - All status update packets currently in the queue are dropped as outdated.
  - A fresh status update packet is sampled and transmitted via a URLLC link (expensive) by the transmitter.
- URLLC link has success probability 1



## Zero-wait

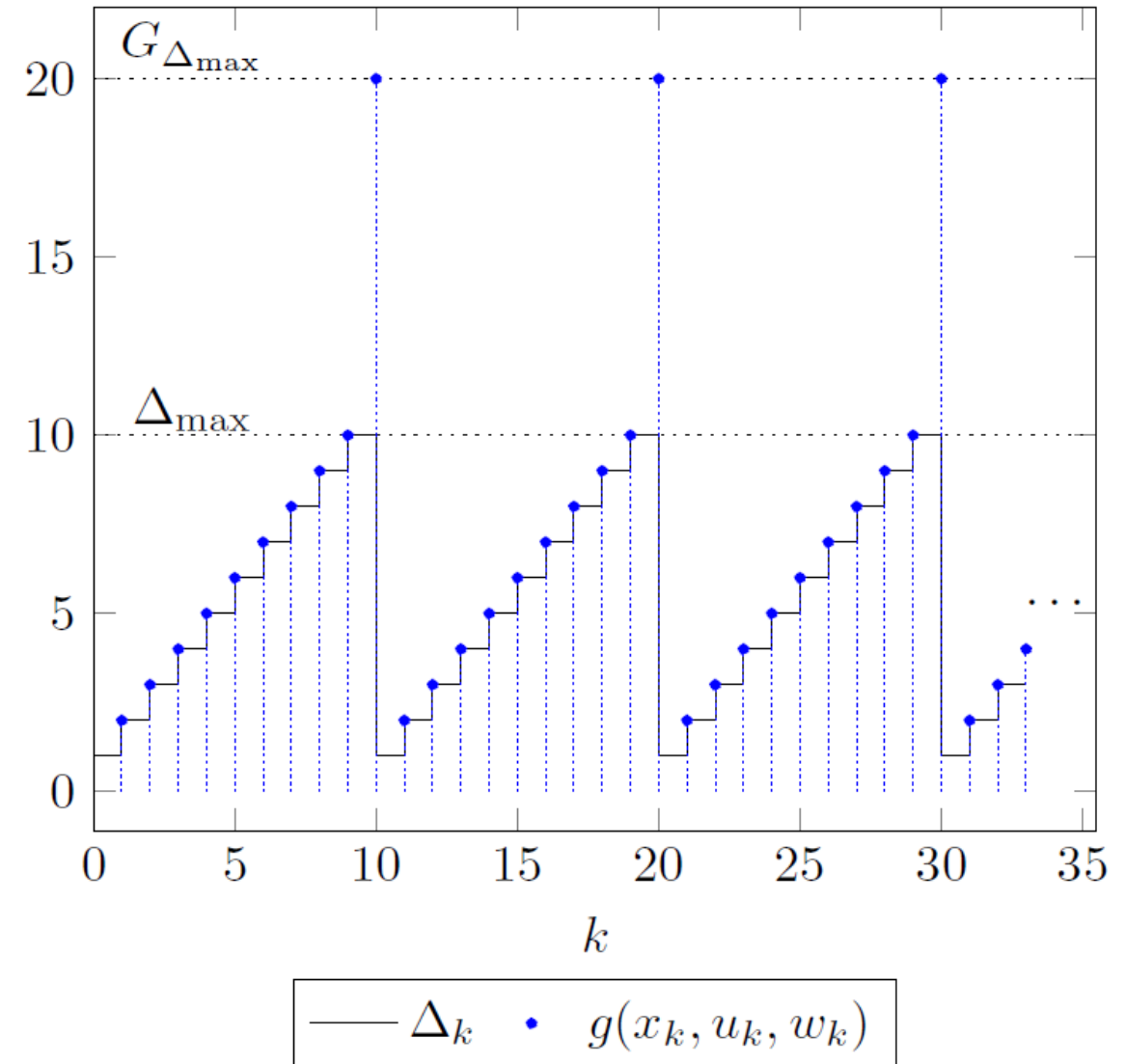
- To achieve zero-waiting in the queue both status update and application packets must have been served prior to generating a new status update.
- The arrival rate of application packets in the queue will determine how often the queue and the server will be idle.

## Max-sampling policy

- Generate a status update at **every** time-slot
- Is meaningful when the application generates packets with a very high rate.
- Queueing delay may result in increased Aol
  - If the application does not generate packets frequently.

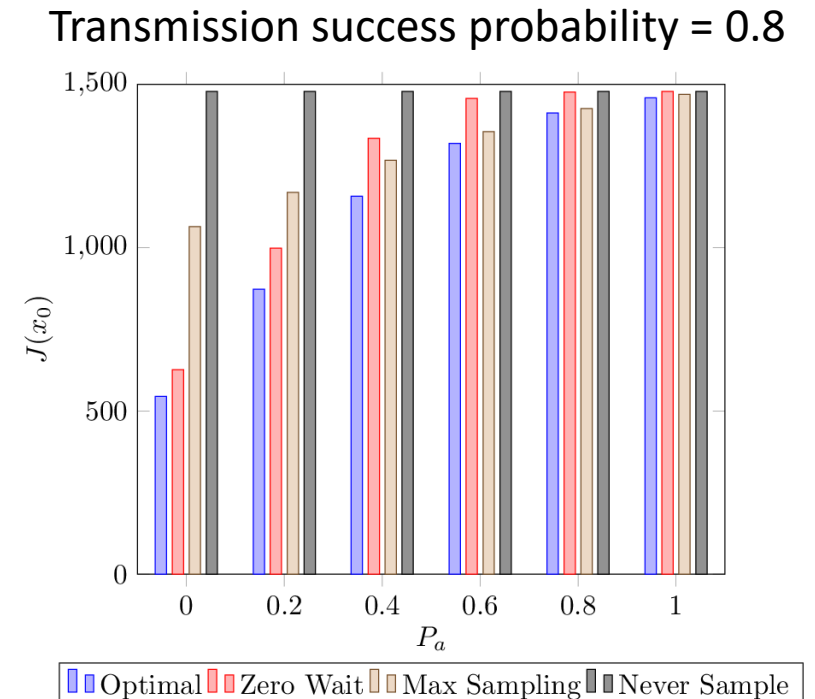
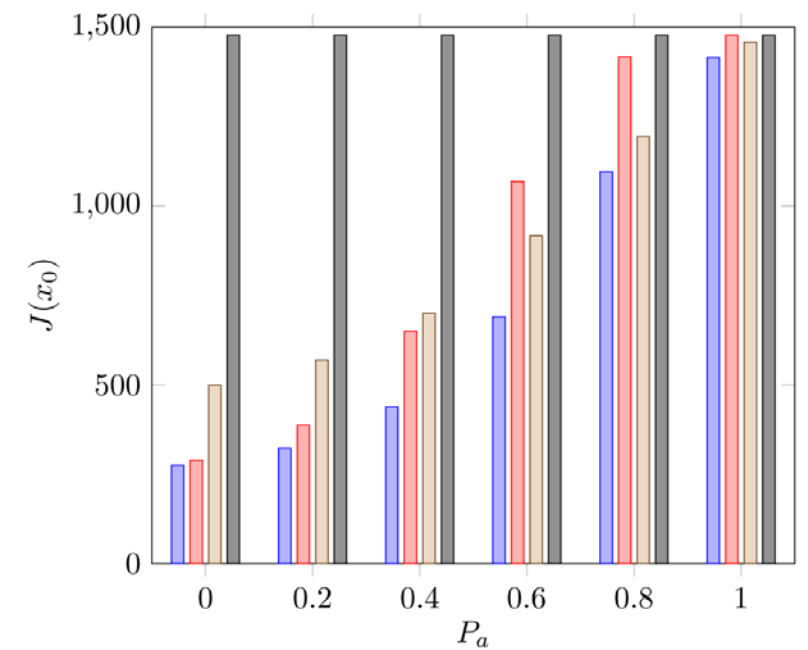
# 3<sup>rd</sup> heuristic policy: Never sample policy

- It never samples the stochastic process unless the upper bound on Aol is violated.
- Subsequently it uses the URLLC mechanism to send a status update.
- This is the worst-performing policy.
  - Complete lack of control over the status update generation process
  - Used to indicate the worst possible performance



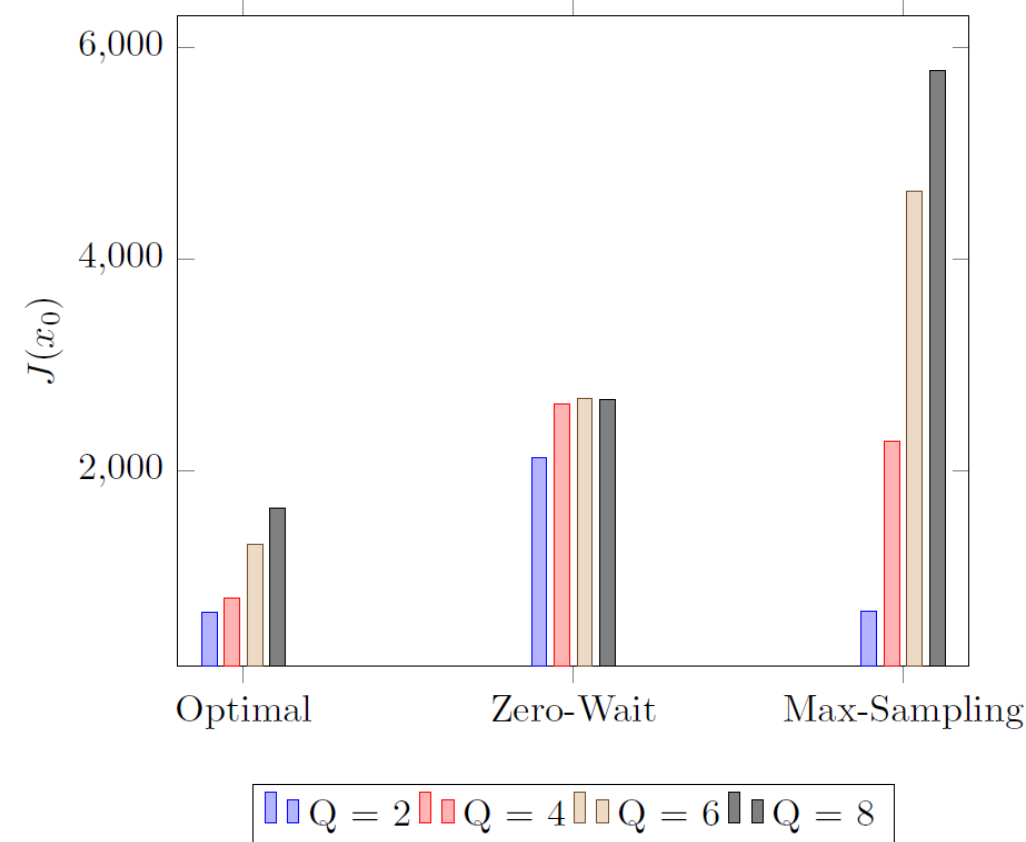
# Results

- System Parameters
  - Queue Size = 4
  - Aol Threshold = 10
  - Max. Retransmission Number = 4
  - URLLC channel cost = 100
  - Discount factor = 0.99
- For low values of  $P_a$  the performance of the zero-wait policy is close to the optimal.
- For higher values of  $P_a$  the performance of the max-sampling policy is close to the optimal policy.
- The never-sample policy has the same performance irrespective of the  $P_a$  value since it only uses the URLLC channel when the Aol threshold is violated.



# Effect of queue size on cost

- System Parameters
  - Queue Size = 2, 4, 6 and 8
  - URLLC channel cost = 1000
  - Transmission success probability = 0.8
  - Application packets arrival probability = 0.4
- The cost for the optimal policy increases with the queue size
  - Increased queueing delays for status updates.
- The max-sampling policy cost increases with queue size
  - Large number of status updates stored in the queue result in increased queueing delay.
- The zero-wait policy is remains unaffected by increasing the size of the queue beyond 4.
  - Increasing the queue size beyond a certain value has no effect on the probability distribution of the number of application packets in the queue.



# Part III Aol-Oriented Multiuser Scheduling

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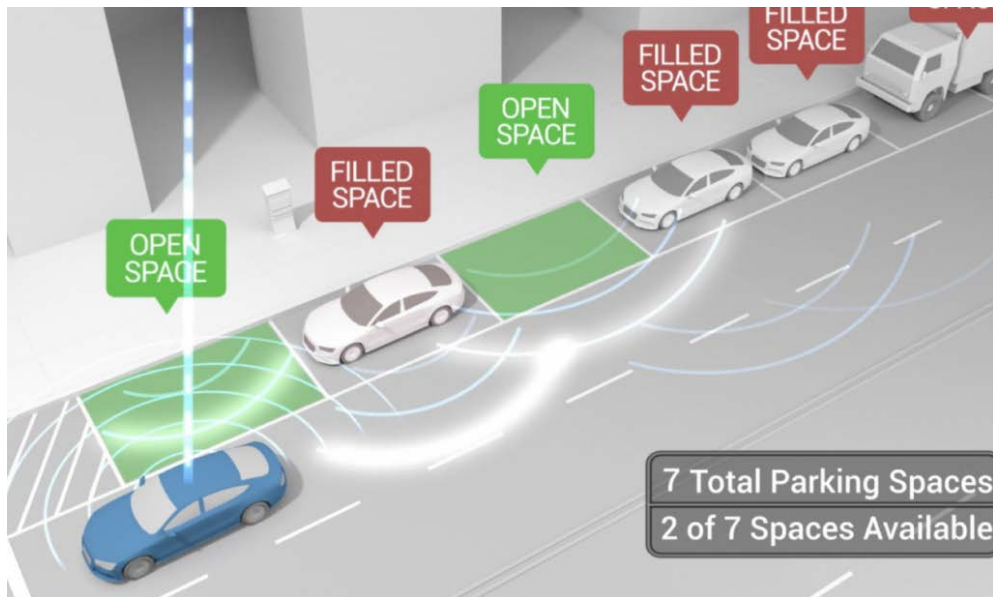
# Multuser Scheduling

- One of the most fundamental problems in wireless networks due to the **shared medium**
- Many work on designing and analyzing **throughput- or delay-oriented multiple scheduling** schemes
- May **no longer optimal** when the Aol is concerned

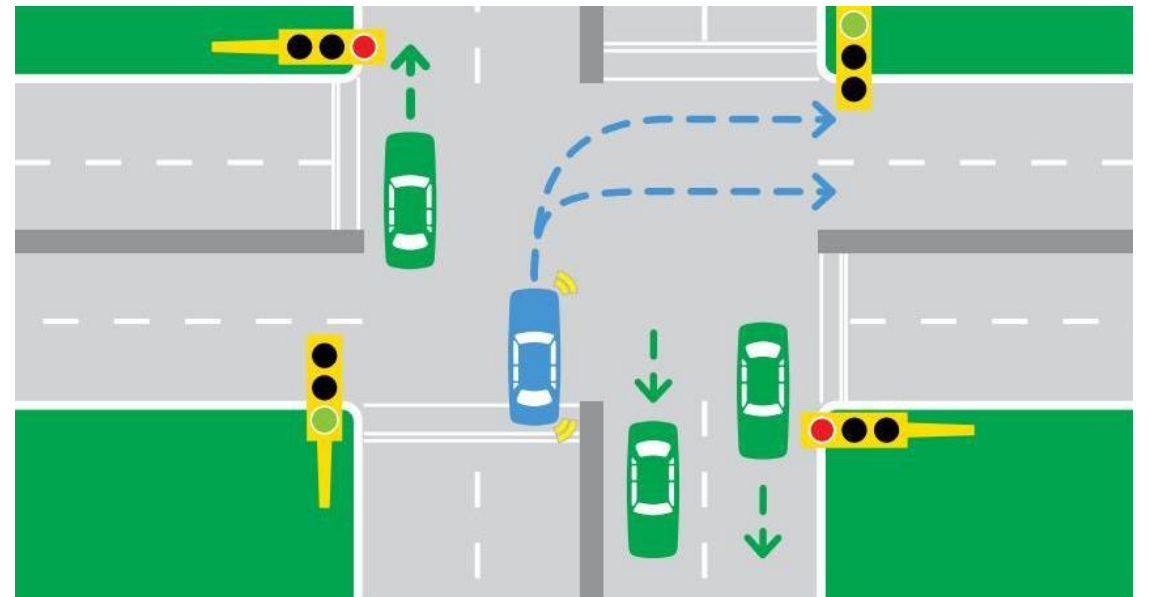
How to schedule the status updates of multiple users to minimize the expected (weighted) sum Aol of the whole system?

# Status update arrival model

- Generate-at-will

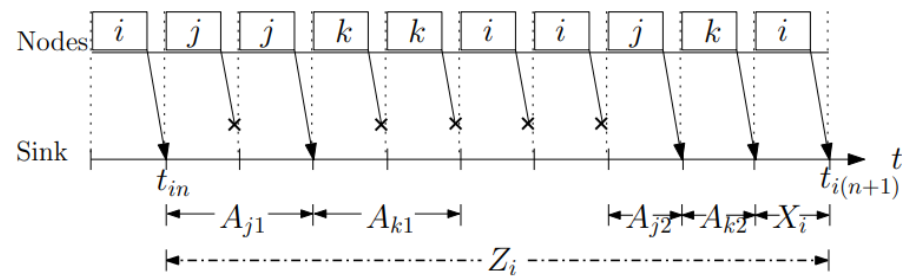


- Random (stochastic) Arrival



# Homogenous Networks

- The channel error probability (when no collision) of all links are **identical**
- Generate-at-will model
- Once scheduled, the node can transmit the its status update for  $S$  times



**Lemma 2.** For a homogeneous network of  $M$  nodes that take turns to transmit their state to a sink and get feedback on whether an update of state occurred, the network's AoI is minimized by allowing a node to keep transmitting packets during its scheduled turn until an update by the node occurs. Specifically, it is minimized in the limit as  $S \rightarrow \infty$ .

- Always schedule the node with **maximum instantaneous age**
- **Maximum-age-first (MAF) policy**
- **MAF is age-optimal** in homogenous networks

# Heterogeneous Networks (1)

- A base station (BS) sending time-sensitive information to  $M$  clients (**downlink**)
- Time is slotted, with  $T$  consecutive slots forming a **frame**, indexed by  $k$ . Let  $n \in \{1, \dots, T\}$  be the index of the slot within a frame
- At the beginning of every **frame**, the BS generates **one packet** per client  $i$
- In each frame,  **$T$  out of  $M$  clients** are scheduled to transmit
- The packet is successfully delivered to client  $i$  with probability  $p_i \in (0, 1]$

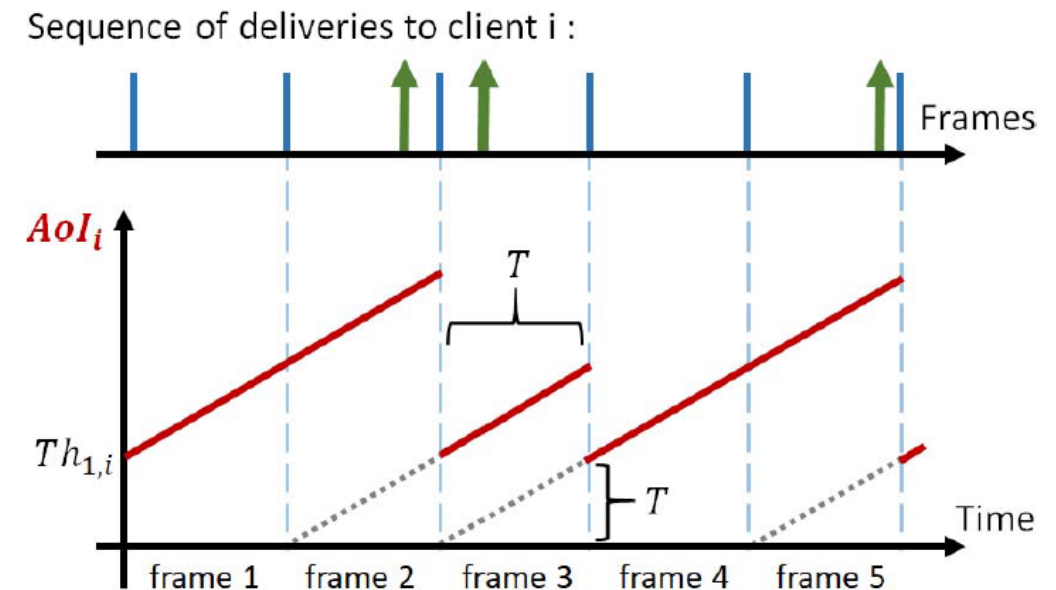


Fig. 2. On the top, a sample sequence of deliveries to client  $i$  during five frames. The upward arrows represent the times of packet deliveries. On the bottom, the associated evolution of the  $AoI_i$ .

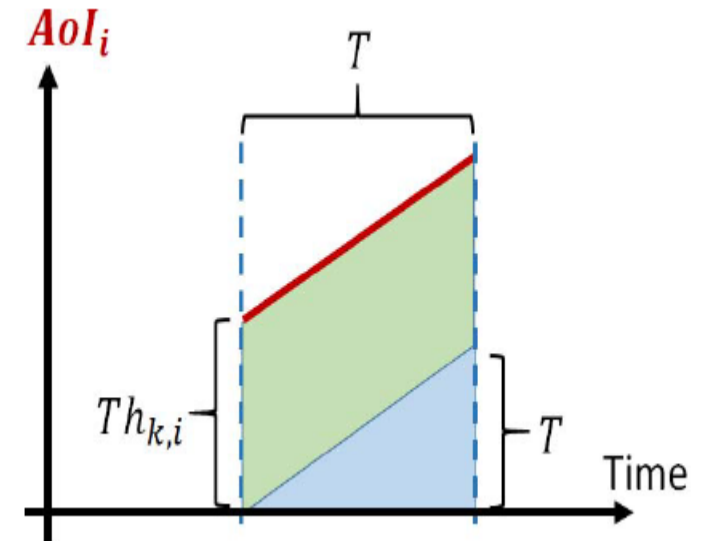
# Heterogeneous Networks (2)

- **Non-anticipative** scheduling policy: policies that **do not use future knowledge** in selecting clients, denoted by  $\Pi$
- $\pi \in \Pi$  be an arbitrary admissible policy
- In a slot  $(k, n)$ , policy  $\pi$  can either idle or select a client with an undelivered packet.
- Performance measure---Expected Weighted Sum AoI

$$\begin{aligned} \text{EWSAoI} &= \frac{1}{KTM} \mathbb{E} \left[ \sum_{k=1}^K \sum_{i=1}^M \alpha_i \left( \frac{T^2}{2} + T^2 h_{k,i} \right) \mid \vec{h}_1 \right] \\ &= \frac{T}{2M} \sum_{i=1}^M \alpha_i + \frac{T}{KM} \mathbb{E} \left[ \sum_{k=1}^K \sum_{i=1}^M \alpha_i h_{k,i} \mid \vec{h}_1 \right], \quad (2) \end{aligned}$$

the positive real value that denotes the client's weight

the number of frames since the last delivery to client  $i$



Area under  $AoI_i$  during any frame  $k$  in terms of  $h_{k,i}$  and  $T$ .

# Heterogeneous Networks (3)

$$\begin{aligned} \text{EWSAoI} &= \frac{1}{KTM} \mathbb{E} \left[ \sum_{k=1}^K \sum_{i=1}^M \alpha_i \left( \frac{T^2}{2} + T^2 h_{k,i} \right) \mid \vec{h}_1 \right] \\ &= \frac{T}{2M} \sum_{i=1}^M \alpha_i + \frac{T}{KM} \mathbb{E} \left[ \sum_{k=1}^K \sum_{i=1}^M \alpha_i h_{k,i} \mid \vec{h}_1 \right], \quad (2) \end{aligned}$$

- Performance formulation (EWSAoI minimization)

$$\min_{\pi \in \Pi} \mathbb{E} [J_K^\pi], \text{ where } J_K^\pi = \frac{1}{KM} \sum_{k=1}^K \sum_{i=1}^M \alpha_i h_{k,i}^\pi,$$

- One possible approach is to use **Dynamic Programming**
- **Computationally demanding**, especially for networks with a large number of clients  $\rightarrow$  low-complexity policy

# Heterogeneous Networks (4)

- **Greedy policy**: in each slot  $(k, n)$ , schedules a transmission to the client with **highest value of  $h_{k,i}$  that has an undelivered packet**

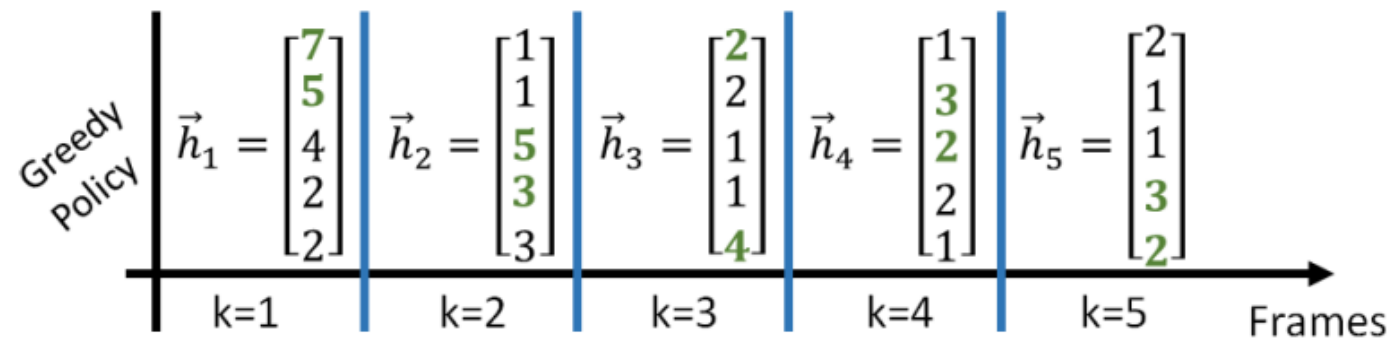


Fig. 4. Evolution of  $\vec{h}_k$  when the Greedy policy is employed in a network with  $M = 5$  clients,  $T = 2$  slots per frame, error-free channels,  $p_i = 1, \forall i$ , and  $\vec{h}_1 = [7 \ 5 \ 4 \ 2 \ 2]^T$ . In each frame, the Greedy policy transmits packets of two clients. The elements of  $\vec{h}_k$  associated with the clients that received a packet during frame  $k$  are depicted in bold green.

# Heterogeneous Networks (5)

- **Properties of Greedy policy**

- It switches scheduling decisions only after a successful packet delivery
- Greedy follows a **Round Robin**-like pattern
- Consider a **symmetric** network with channel reliabilities  $p_i = p \in (0, 1]$  and weights  $\alpha_i = \alpha > 0, \forall i$ . Among the class of admissible policies  $\Pi$ , the **Greedy policy attains the minimum expected sum Aol**



# Heterogeneous Networks (6)

- **Stationary Randomized Policy**

*Randomized policy selects in each slot  $(k, n)$  client  $i$  with probability  $\beta_i / \sum_{j=1}^M \beta_j$ , for every client  $i$  and for positive fixed values of  $\{\beta_i\}_{i=1}^M$ . The BS transmits the packet if the selected client has an undelivered packet and idles otherwise.*

- this policy uses **no information** from current or past states of the network.

# Heterogeneous Networks (7)

- **Max-Weight Policy**

- Derived based on Lyapunov Optimization

*Max-Weight policy schedules in each slot  $(k, n)$  a transmission to the client with highest value of  $p_i \alpha_i h_{k,i} (h_{k,i} + 2)$  that has an undelivered packet, with ties being broken arbitrarily.*

- Observe that when  $\alpha_i = \alpha$  and  $p_i = p$ , prioritizing according to  $p_i \alpha_i h_{k,i} (h_{k,i} + 2)$ , i.e. Max-Weight is identical to Greedy.
- **Max-Weight is Aol-optimal** for symmetric networks.

# Heterogeneous Networks (8)

- **Whittle's Index policy**

- the optimal solution to a **relaxation** of the Restless Multi-Armed Bandit (RMAB) problem
- First step is to realize that each client in the Aol problem evolves as a restless bandit
- Next step is to consider the **relaxed version** of the RMAB problem, called the Decoupled Model, in which **clients are examined separately**.
- The Decoupled Model associated with each client  $i$  adheres to the network model with  $M = 1$ , except for the addition of a **service charge  $C$** .
- The service charge is a fixed cost per transmission  $C > 0$  that is incurred by the network every time the BS transmits a packet.
- The final step is to prove that the **Aol problem is indexable** and derive the Whittle's Index policy

# Heterogeneous Networks (9)

- Whittle's Index policy

*Whittle's Index policy schedules in each slot  $(k, n)$  a transmission to the client with highest value of*

$$C_i(h_{k,i}) = p_i \alpha_i h_i \left[ h_i + \frac{1 + (1 - p_i)^T}{1 - (1 - p_i)^T} \right],$$

*that has an undelivered packet,*

- Look **similar** to the Max-Weight policy
- Both are equivalent to the Greedy policy in symmetric networks
- Both the Whittle's Index and Max-Weight policies have strong performances

# Heterogeneous Networks (10)

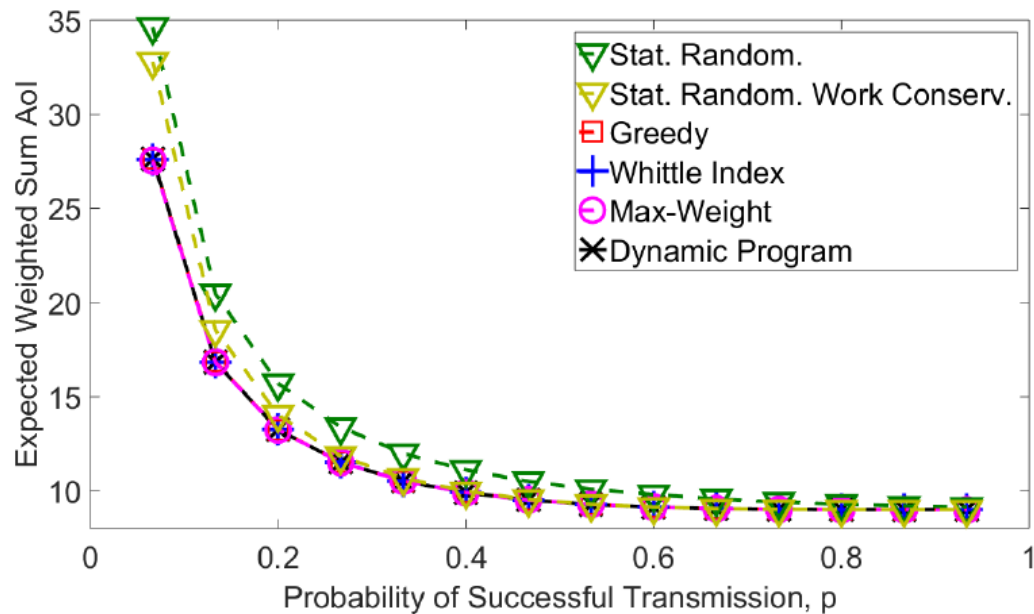


Fig. 5. Two-user symmetric network with  $T = 6, K = 150, \alpha_i = 1, p_i = p, \forall i$ . The simulation result for each policy and for each value of  $p$  is an average over 1,000 runs.

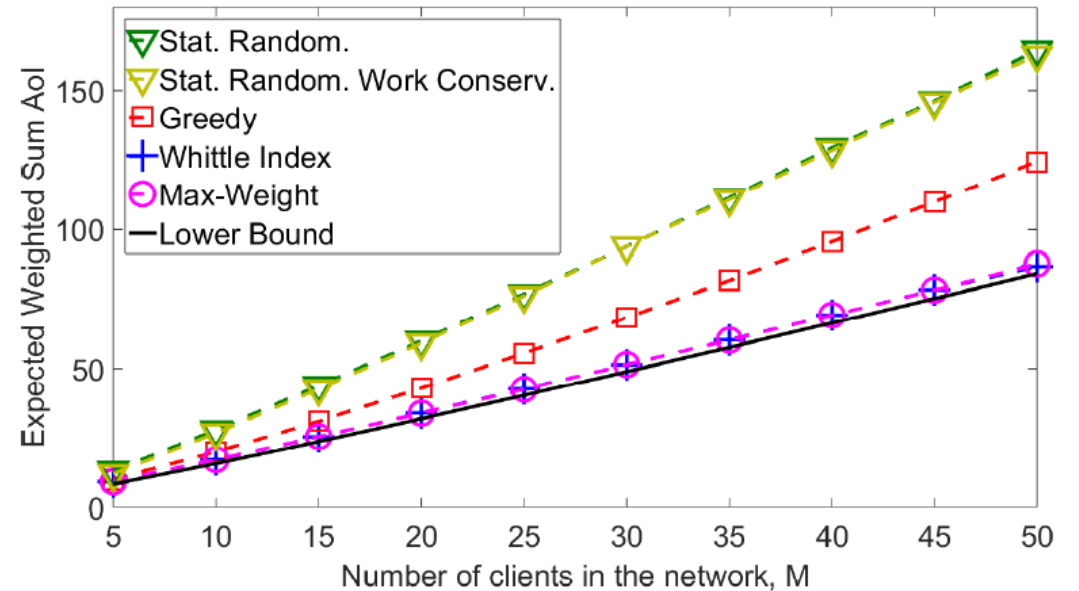
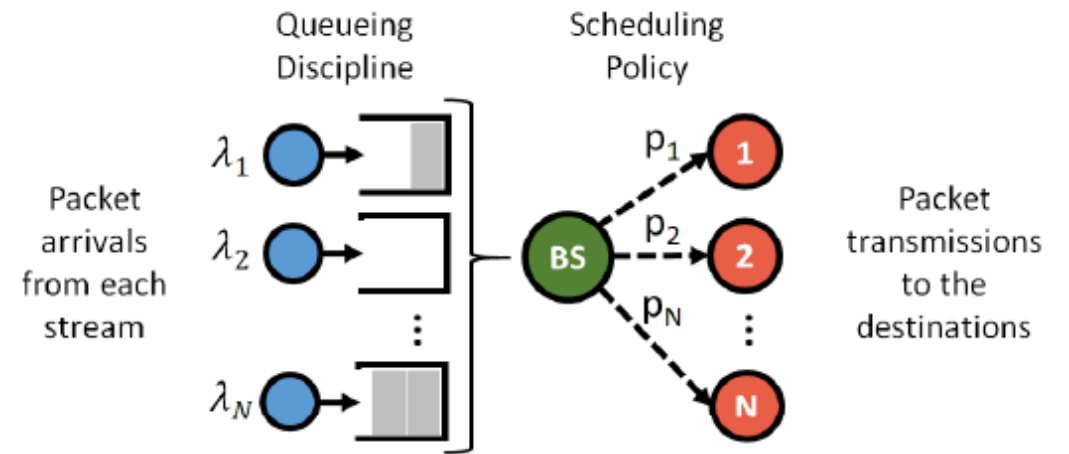


Fig. 7. Network with  $T = 3, K = 50,000, \alpha_i = 1, p_i = i/M, \forall i$ . The simulation result for each policy and for each value of  $M$  is an average over 10 runs.

# Random arrival case (1)

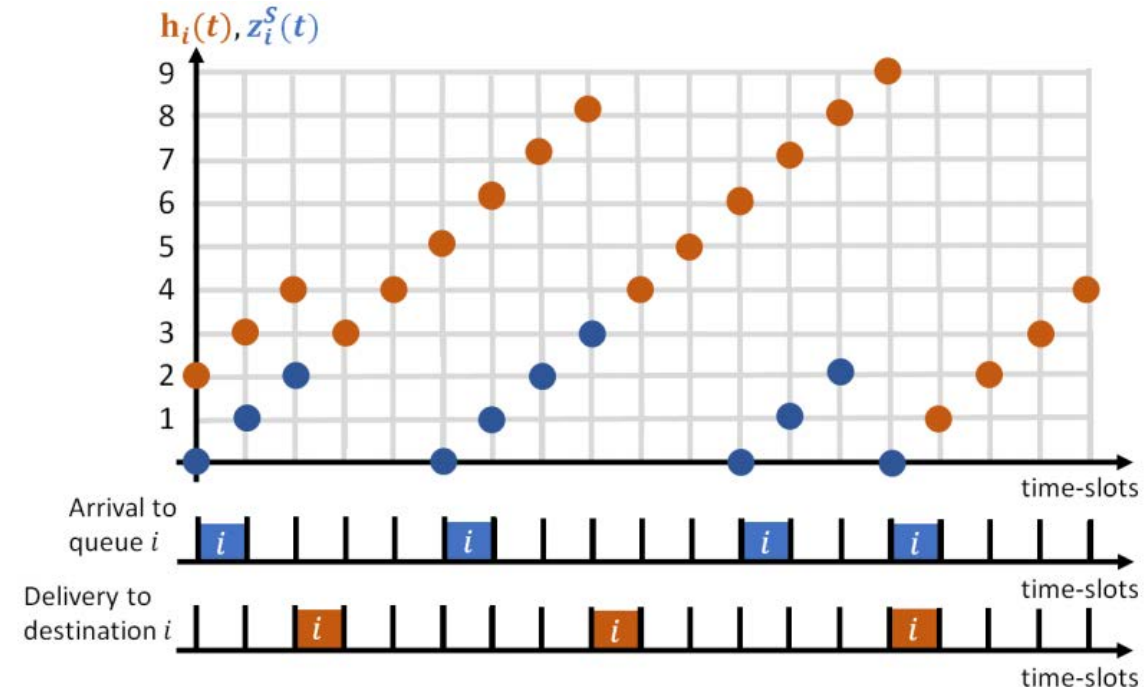
- One BS, N users, **downlink**
- a new packet to user  $i \in \{1, 2, \dots, N\}$  arrives to the system with **probability**  $\lambda_i \in (0, 1]$ ,
- Three queueing disciplines
- Successful transmission probability  $p_i$
- The BS can transmit **at most one packet** at any given timeslot  $t$



- Three Queueing Discipline:
- (1) First In First Out (FIFO) queues
  - (2) Single Packet queues
  - (3) No queues

# Random arrival case (2)

- Key difference from the generate-at-will model: *The system time of the packets should be jointly considered together with the Aol*
- E.g., a packet that has waited for a long time may not worth a transmission
- $h_i$  denotes Aol,  $z_i$  denotes the system time



# Random arrival case (3)

- To minimize the Expected Weighted Sum AoI (EWSAoI)

$$\mathbb{E}[J^\pi] = \lim_{T \rightarrow \infty} \frac{1}{TN} \sum_{t=1}^T \sum_{i=1}^N w_i \mathbb{E}[h_i^\pi(t)]$$

- **Max-Weight** policies:

*in each slot  $t$ , the stream  $i$  with a HoL packet and the highest value of  $\beta_i p_i (h_i(t) - z_i(t))$ ,*

- $\beta_i$  is a positive **hyperparameter** that can be used to tune the Max-Weight policy to different network configurations and queueing disciplines.
- the difference  $h_i(t) - z_i(t)$  represents the **AoI reduction** accrued from a successful packet delivery to destination  $i$

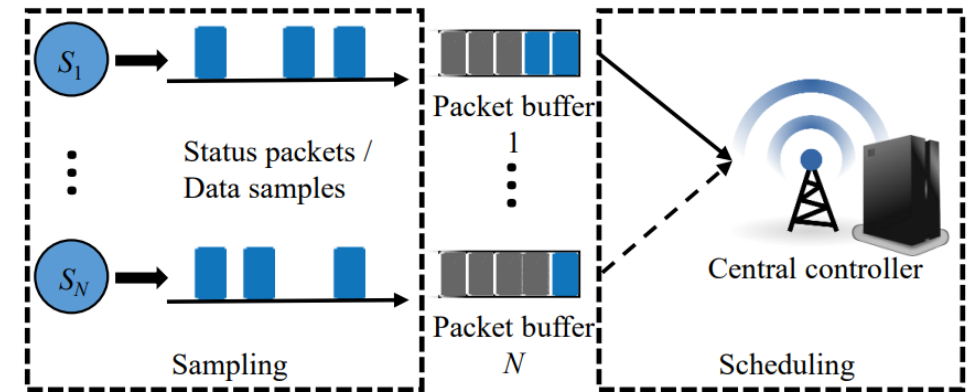


# Random arrival case in the uplink

- More **complicated**: the BS/monitor may not know whether new status updates arrive at end nodes ( $z_i(t)$  is not exactly known)
- The BS/monitor needs to make a scheduling decision under **partially observable system states** (Instantaneous Aol is known)
- [Gong'20] addressed the EWSAol minimization by formulating it as a **partially observable Markov decision process** (POMDP)
- **Belief state** characterizes the fully observable Aol and the partially observable status update arrivals of end nodes at the monitor

# Joint Design of Sampling and Scheduling

- Sampling cost (e.g., energy consumption) is considered
- Generate-at-will is no longer age-optimal
- When to sample and when to transmit need to be jointly optimized



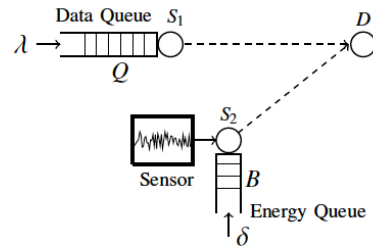
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[Zhou'19] Bo Zhou and Walid Saad, "Joint Status Sampling and Updating for Minimizing Age of Information in the Internet of Things", IEEE Transactions on Communications, vol. 67, no. 11, pp. 7468-7482, Nov. 2019.

[Jiang'19] Z. Jiang, S. Zhou, Z. Niu and C. Yu, "A Unified Sampling and Scheduling Approach for Status Update in Multiaccess Wireless Networks," IEEE INFOCOM 2019 - IEEE Conference on Computer Communications, Paris, France, 2019.

[Fountoulakis'20] E. Fountoulakis, N. Pappas, M. Codreanu, A. Ephremides, "Optimal Sampling Cost in Wireless Networks with Age of Information Constraints", IEEE INFOCOM - 3rd AoI Workshop 2020.

# Hybrid Systems



- ▶ Time-slotted multiple access channel (MAC).
- ▶  $S_1$ : connected to a power grid, Bernoulli data arrivals with arrival probability  $\lambda$ , queue backlog  $Q$ .
- ▶  $S_2$ : energy harvesting (EH) sensor, Bernoulli energy arrivals with probability  $\delta$ , battery level (energy queue)  $B$ .
- ▶ Multi-packet reception (MPR) capabilities at the destination node  $D$ .

## Transmission and Re-transmission Policy

### Transmission policy

- ▶  $S_1$ : First In, First Out (FIFO);
- ▶  $S_2$ : always generates a fresh sample before attempted transmission.

### In case of failed transmission

- ▶  $S_1$ : re-transmits the same packet in the next slot;
- ▶  $S_2$ : drops that packet and generates a new sample before the next attempted transmission.

- ▶ Stationary randomized policy: random access decisions based on some stationary probability distribution.
- ▶ Drift-Plus-Penalty policy: Lyapunov optimization, virtual queue, drift-plus-penalty.

## Simulation Results

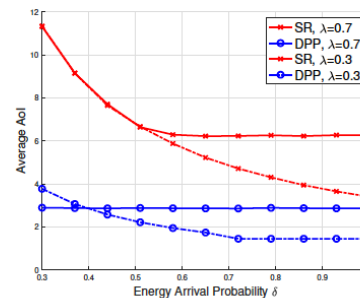


Figure: Weak MPR.

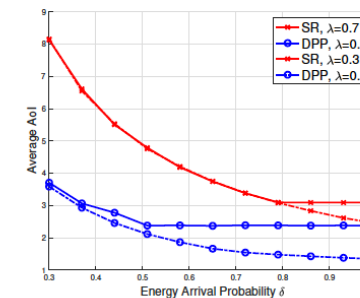


Figure: Strong MPR.

- DPP policy performs significantly better than the SR policy, especially when the destination node has weak MPR capabilities.

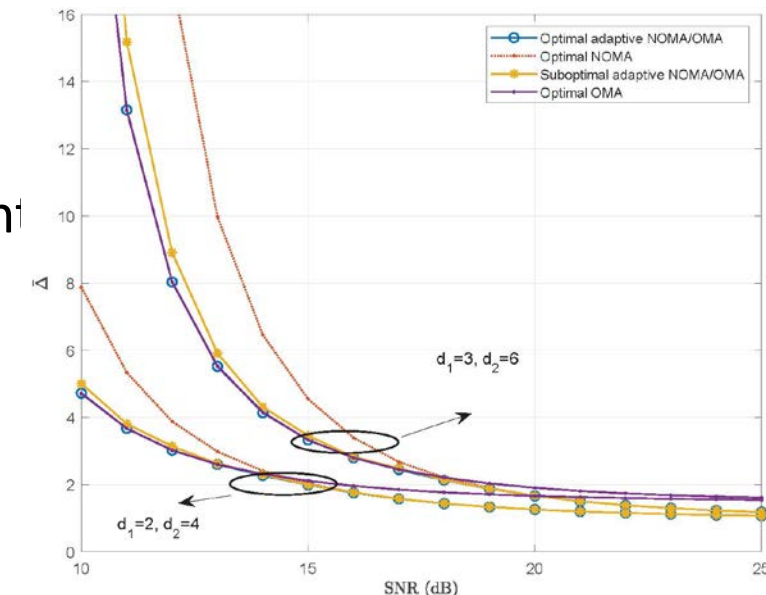
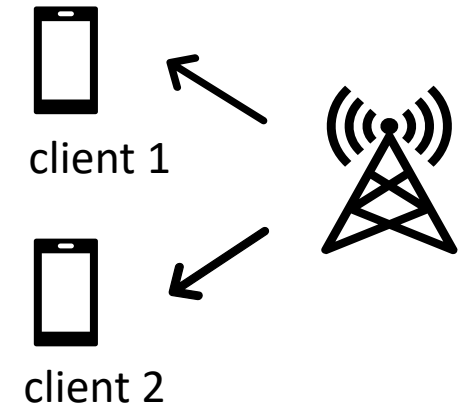
[Chen-Pappas'20] Z. Chen, N. Pappas, E. Björnson, E. G. Larsson, "Optimizing Information Freshness in a Multiple Access Channel with Heterogeneous Devices", arXiv:1910.05144v2, July 2020.

[Wang'20] Q. Wang, H. Chen, Y. Gu, Y. Li, and B. Vucetic, "Minimizing the age of information of cognitive radio-based IoT systems under a collision constraint," <https://arxiv.org/pdf/2001.02482.pdf>, Jan. 2020.

[Moltafet'20] M. Moltafet, M. Leinonen, M. Codreanu, N. Pappas, "Power Minimization for Age of Information Constrained Dynamic Control in Wireless Sensor Networks", arXiv:2007.05364, July 2020.

# Scheduling 1+ users per timeslot (1)

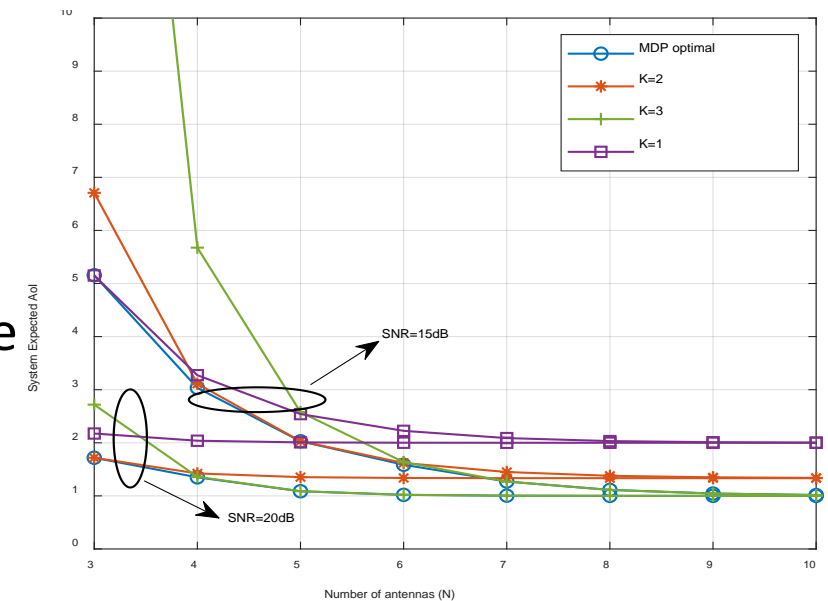
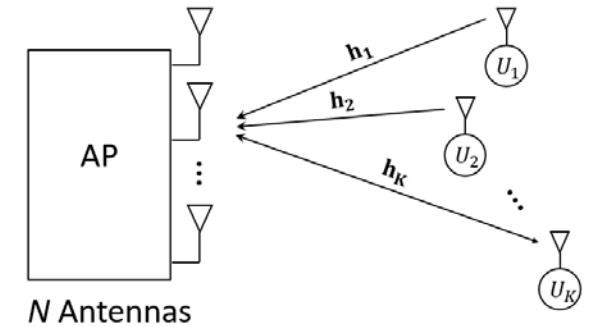
- Users can update **more frequently** → lower average Aol
- [Wang'20] proposed to apply **adaptive NOMA/OMA**
- A base station (BS) conducting timely transmission to **two clients** in a slotted manner via adaptive NOMA/OMA.
- generate-at-will model was considered
- In **OMA**, the BS decides which client to conduct transmission.
- In **NOMA**, the BS should determine power allocated to each client
- **Decision to make**: power allocation between the two users
- Formulated as a Markov decision process (MDP) problem
- Low-complexity **max-weight policy** was derived



[Wang'20] Q. Wang, H. Chen, Y. Li, and B. Vucetic, "Minimizing age of information via hybrid NOMA/OMA," presented at ISIT 2020, available: <https://arxiv.org/abs/2001.04042>

# Scheduling 1+ users per timeslot (2)

- [Chen'20] proposed to use **multi-antenna techniques**
- Uplink, generate-at-will
- K users, one AP with N antennas
- Can schedule **N users in each time slot**
- An inherent **trade-off** exists in multiuser MIMO systems: *scheduling more users to transmit in the same time slot will lead to a higher transmission error probability for each scheduled user.*
- Decision to make: which **user or a group of users** to schedule in each time slot
- Formulated as an MDP problem
- **Max-weight policy** was also studied



# Part IV Aol-Oriented Random Access

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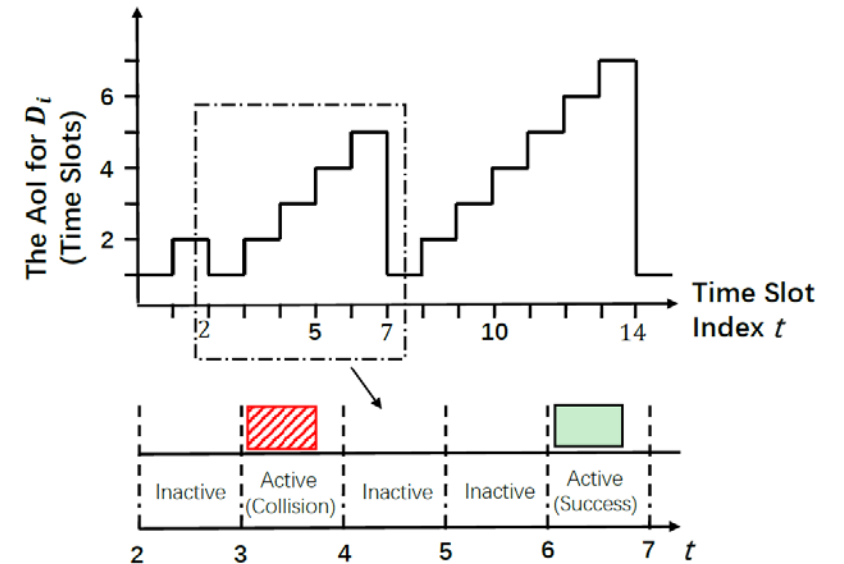
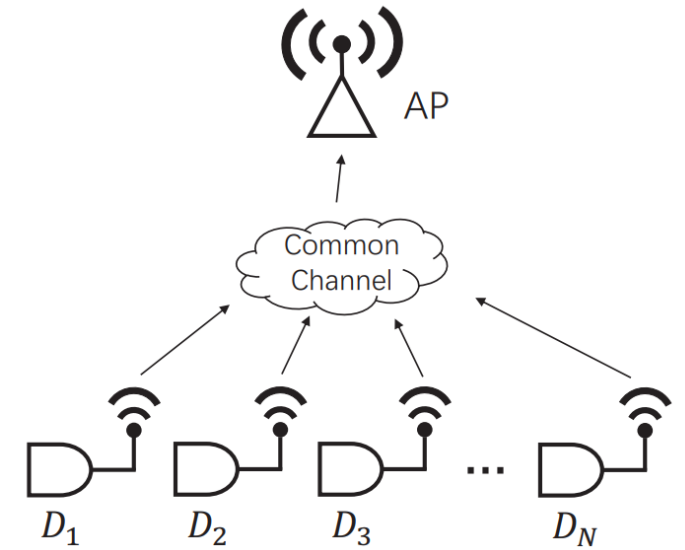
# Why random access?

- Massive connectivity is coming → Ericsson foresaw that by 2021, there will be around **28 billion** IoT devices
- Centralized scheduling → overhead is **excessive** for a large-scale network
- Decentralized scheduling (**random access**) is more preferred

How to schedule the status updates of massive IoT devices to achieve a low network-wide AoI?

# Generic Model

- Consider an uplink network consisting of **one access point (AP)** and  **$N$  IoT devices** ( $N$  can be large)
- **Time is divided into slots of equal durations**
- the transmission of each status update packet takes one time slot





# Age-independent random access (AIRA)

- **Slotted ALOHA-like** random access
- The active probability for each IoT device  $p'$  is **independent** of its instantaneous AoI
- Generate-at-will model
- Let  $q$  denote the **successful status update probability** when an IoT device becomes active

$$q = (1 - p')^{N-1}$$

- Average AoI (collision channel)

$$\bar{\Delta}' = \frac{1}{p'q} = \frac{1}{p'(1 - p')^{N-1}}.$$

# Age-Dependent Random Access (ADRA)

- [Chen-Gu'20] proposed **age-dependent** random access
- **Intuition**: those devices with relatively smaller Aol should access the channel with a lower probability such that other devices with larger Aol can achieve a higher success probability to update their statuses by encountering less collisions
- For simplicity, we consider a **threshold-based ADRA** protocol: if the instantaneous Aol is no less than a **threshold  $\delta$** , the IoT device becomes activate with a **fixed probability of  $p$** . Otherwise, the IoT device will stay **inactive with probability  $1$**

# Average Aol Analysis for ADRA (1)

- In the AIRA policy, the successful status update probability  $q$  of one device is **independent** of the instantaneous Aol of all other devices
- In the proposed ADRA protocol, the active probability of each IoT device (i.e.,  $p$ ) **depends on** its instantaneous Aol.
  - ➔  $q$  depends on the instantaneous Aol of all IoT devices and thus the Aol evolutions of all devices tangle together
- **Multi-dimension** Markov Chain (MC) is a possible solution, however the computational complexity is very high

# Average Aol Analysis for ADRA (2)

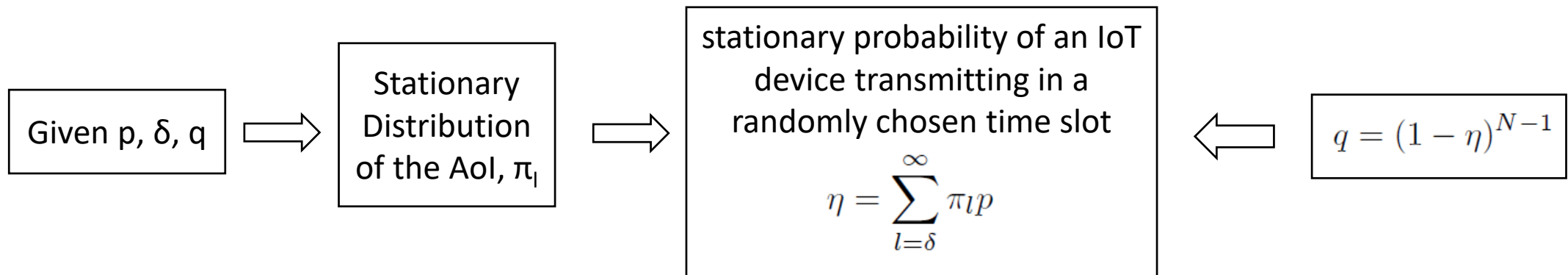
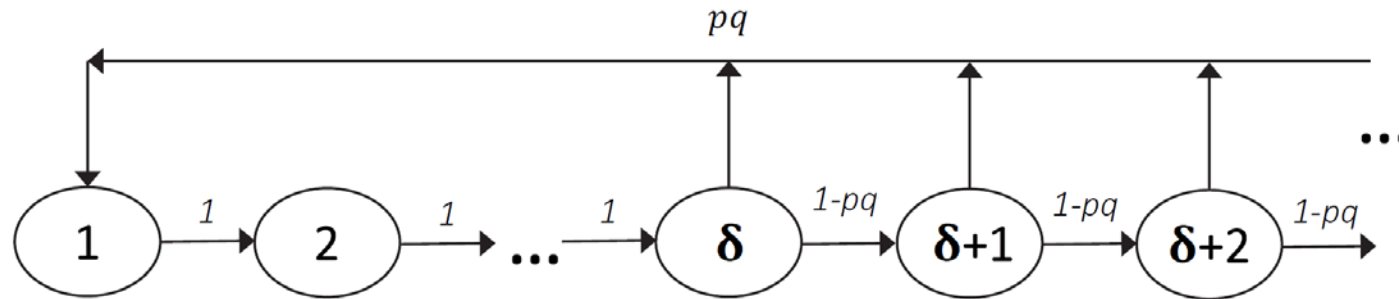
- **Key approximation** to decouple the tangled evolution of the Aol for all IoT devices:

*The **successful probability  $q$**  for all IoT devices is a constant when they decide to transmit a status update.*

- The value of  $q$  is **independent** of the instantaneous Aol of all other IoT devices, but still **a function of the age threshold  $\delta$  and the active probability  $p$** .
- The accuracy of such approximation has been proved in the literature

# Average Aol Analysis for ADRA (3)

- All devices follow an identical state transition process, which can be described by a **Discrete-Time Markov Chain**



# ADRA with random arrival

- **Local Aol** needs to be considered
- M identical source nodes, collision channel model
- New status update arrives with **probability  $\theta$**  in each time slot and replace the staled one in the buffer
- slotted ALOHA: At every time slot  $k$ , transmitters send their packets immediately upon arrival unless they are “backlogged” after a collision in which case they transmit with a back-off probability
- Rivest’s stabilized slotted ALOHA: all arrivals are regarded as backlogged nodes that transmit with the **backoff probability  $p_b(k)$**
- Let  $c(k) = 1$  denote the event that collision occurred at time  $k$  and  $c(k) = 0$  denote the complementary event

$$p_b(k) = \min\left(1, \frac{1}{n(k)}\right)$$
$$n(k) = \begin{cases} \min\left(n(k-1) + M\theta + (e-2)^{-1}, M\right) & \text{if } c(k) = 1 \\ \min\left(\max(M\theta, n(k-1) + M\theta - 1), M\right) & \text{if } c(k) = 0 \end{cases}$$

- **Conventional slotted ALOHA is asymptotically optimal in terms of the Normalized Expected Weighted Sum Aol (NEWSAol) when  $\theta < 1/(eM)$**
- When  $\theta > 1/(eM)$ , local Aol-aware (**threshold-based**) schemes were designed.

# Performance Comparison

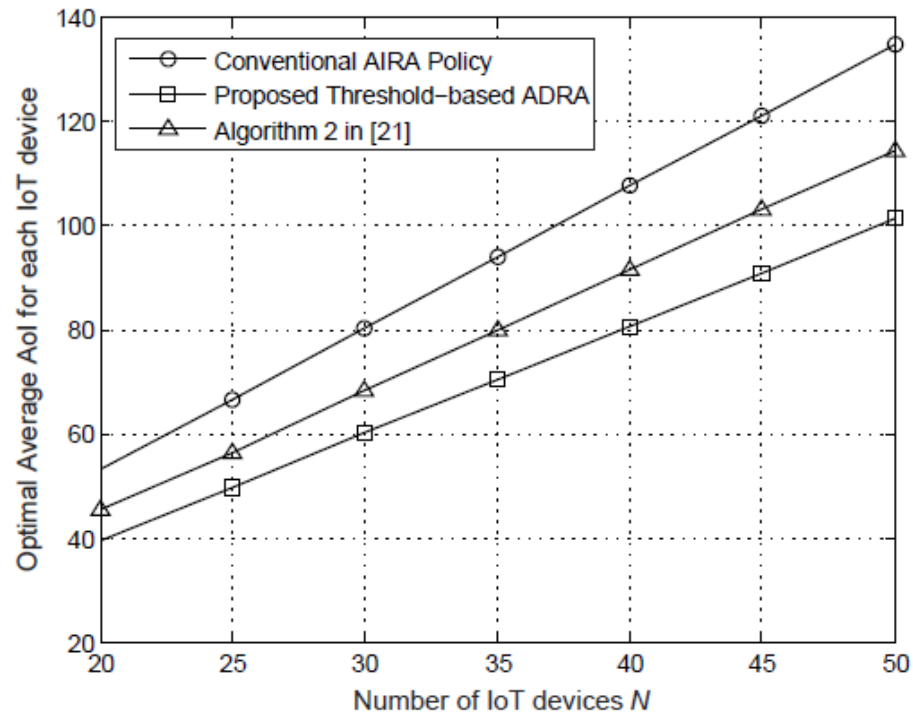


Fig. 5. The average AoI versus the number of IoT devices  $N$  for the proposed threshold-based ADRA, the existing AIRA schemes and Algorithm 2 in [21].

# Threshold-ALOHA

- **Threshold-ALOHA** [Yavascan'20] : sources will wait until their age reaches a certain **threshold  $\Gamma$** , before they turn on their slotted ALOHA mechanism, and only then start to attempt transmission with a fixed **probability  $\tau$**  at each time slot
- Derived time-average expected Age of Information (AoI) attained by this policy, and explored its scaling with **network size,  $n$**
- the **optimal** age threshold and transmission probability are  **$2.2n$  and  $4.69/n$** , respectively
- the optimal AoI scales with the network size as  **$1.4169n$** , which is almost half the minimum AoI achievable using slotted ALOHA, while the loss from the maximum achievable throughput of  $e^{-1}$  remains below **1%**.



# Index-Prioritized Random Access

- Whittle index scheduling policies can achieve near-optimal AoI but require heavy **signalling overhead**
- [Sun-Jiang'20] proposed a contention-based random-access scheme, namely Index-Prioritized Random Access (IPRA)
- Each terminal can calculate **its own index**
- This individual index can be mapped to a **transmission probability** based on a public mapping function which captures the idea that only valuable packets (packets with high index values) are transmitted
- A **single-threshold** function was used in [Sun-Jiang'20]

# Carrier sensing multiple access (CSMA)

- [Wang'19] studied the broadcast Age of Information of CSMA/CA networks
- [Maatouk'20] analyzed the average AoI of CSMA for both generate-at-will and stochastic arrival models, optimized the CSMA scheme and proposed a modification to it by giving **each link the freedom to transition to SLEEP mode**
- [Bedewy'20] **optimized the sleep-wake parameters** for minimizing the weighted-sum peak AoI of the sources, subject to per-source battery lifetime constraints

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[Wang'19] M. Wang and Y. Dong, "Broadcast Age of Information in CSMA/CA Based Wireless Networks," 2019 15th International Wireless Communications & Mobile Computing Conference (IWCMC), Tangier, Morocco, 2019, pp. 1102-1107

[Maatouk'20] A. Maatouk, M. Assaad and A. Ephremides, "On the Age of Information in a CSMA Environment," in IEEE/ACM Transactions on Networking, vol. 28, no. 2, pp. 818-831, April 2020, doi: 10.1109/TNET.2020.2971350.

[Bedewy'20] Ahmed M. Bedewy, Yin Sun, Rahul Singh, and Ness B. Shroff, "Optimizing Information Freshness using Low-Power Status Updates via Sleep-Wake Scheduling", accepted by ACM MobiHoc, 2020.

# Part V Prototype for Experimental Study of Aol-oriented Designs

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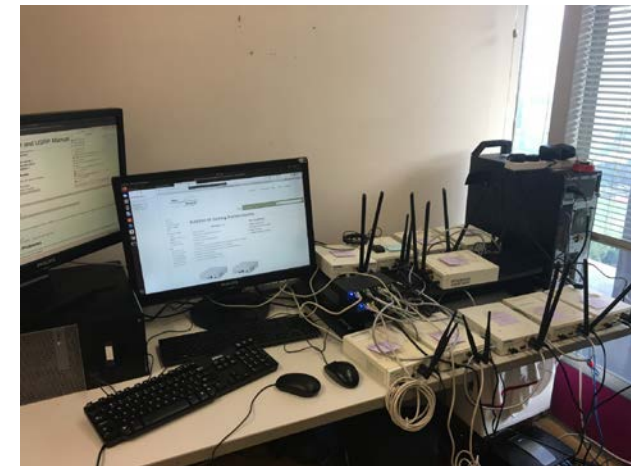
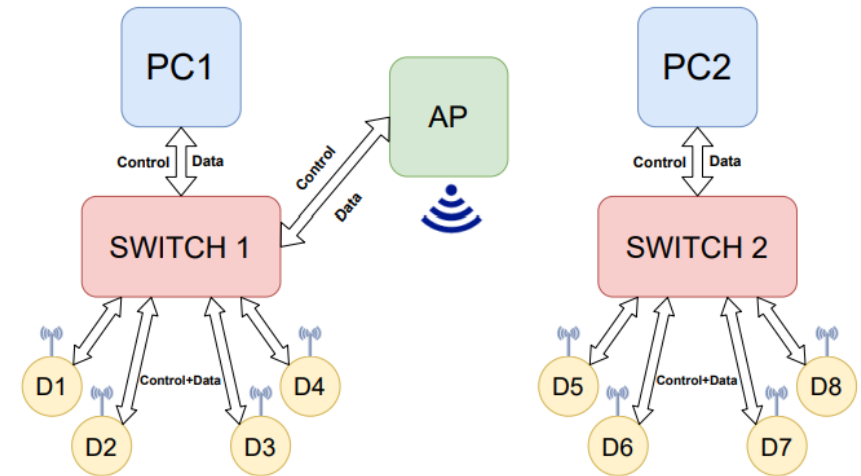
<http://iiotc.ie.cuhk.edu.hk/>

2020.08.09 @ IEEE/CIC ICC 2020 Tutorial



# SDR-based proof-of-concept prototype

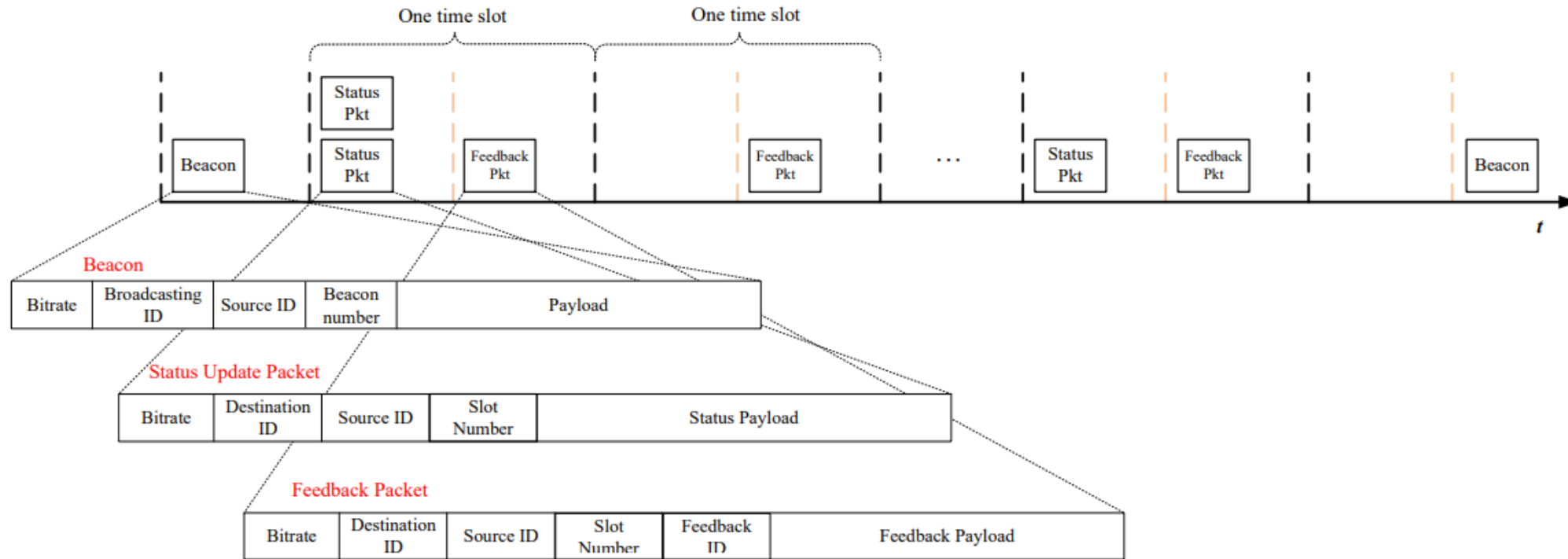
- **Ettus USRP N210's as transceivers**: One USRP N210 serves as the AP and eight USRP N210s serve as end devices
- All of them are connected to two powerful PCs through multiple 1 Gigabit Ethernet cables and two Ethernet switches
- though some USRPs are connected to the same PC, **they use individual local clocks on their motherboards**
- The SBX RF frontboards embedded in USRPs are used to transmit RF signals, **working at 1 GHz with 500 kHz channel bandwidth**.
- We use the **GNURadio platform** to define the signal generation and data processing in our SDR prototype



# SDR-based proof-of-concept prototype

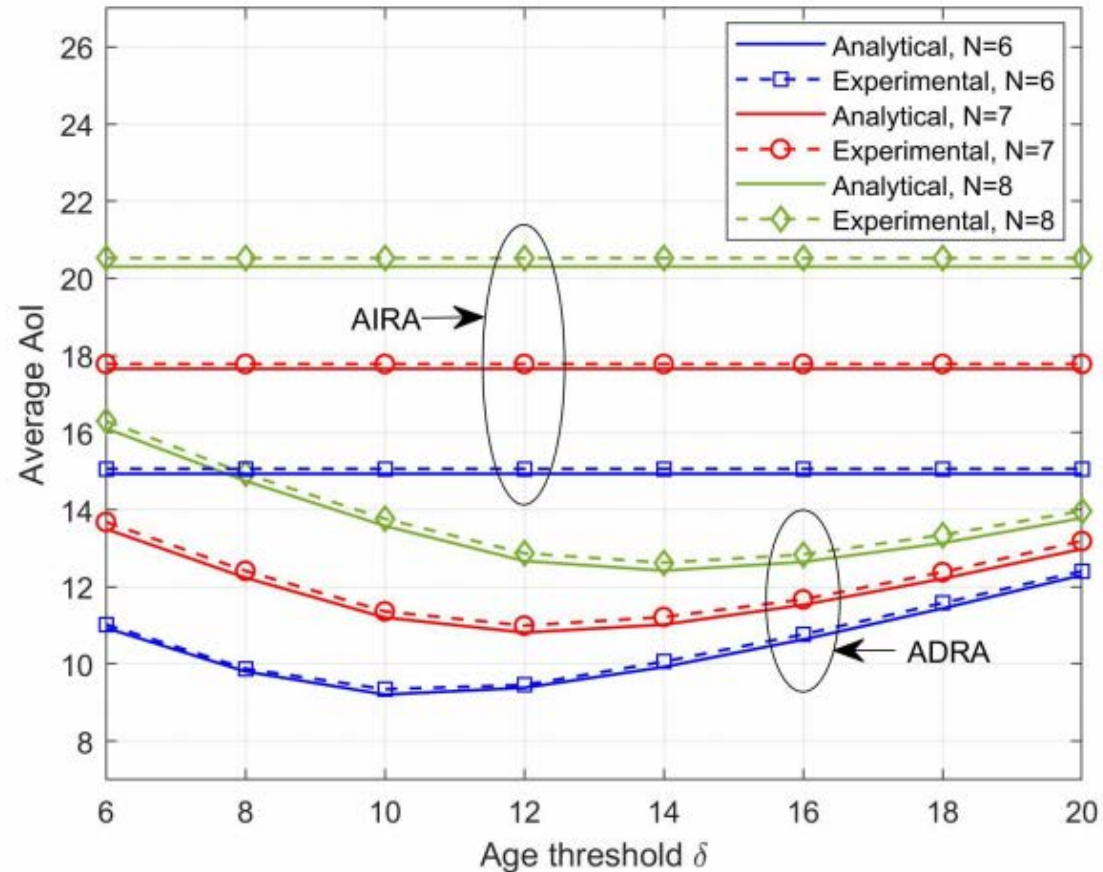
- We implemented the **ADRA protocol** and tested it in office environments
- To mimic the collision channel, we implement the **convolutional code** in the physical layer. **Orthogonal frequency duplex modulation (OFDM)** is used for higher frequency efficiency
- A simple yet effective synchronized transmission scheme for time-slotted transmission
  - (1) **beacon broadcasting**: aims to achieve the time synchronization (the USRP that acts as the AP broadcasts a Beacon, which contains timing information to serve as a time reference, the inter-beacon period is set as 100 time slots in our experiments)
  - Each time slot is further split to two parts for channel access and feedback
  - (2) **channel access**: status update by end devices
  - (3) **feedback**: AP broadcasts a Feedback Packet indicating the successful reception of a Status Update Packet

# SDR-based proof-of-concept prototype



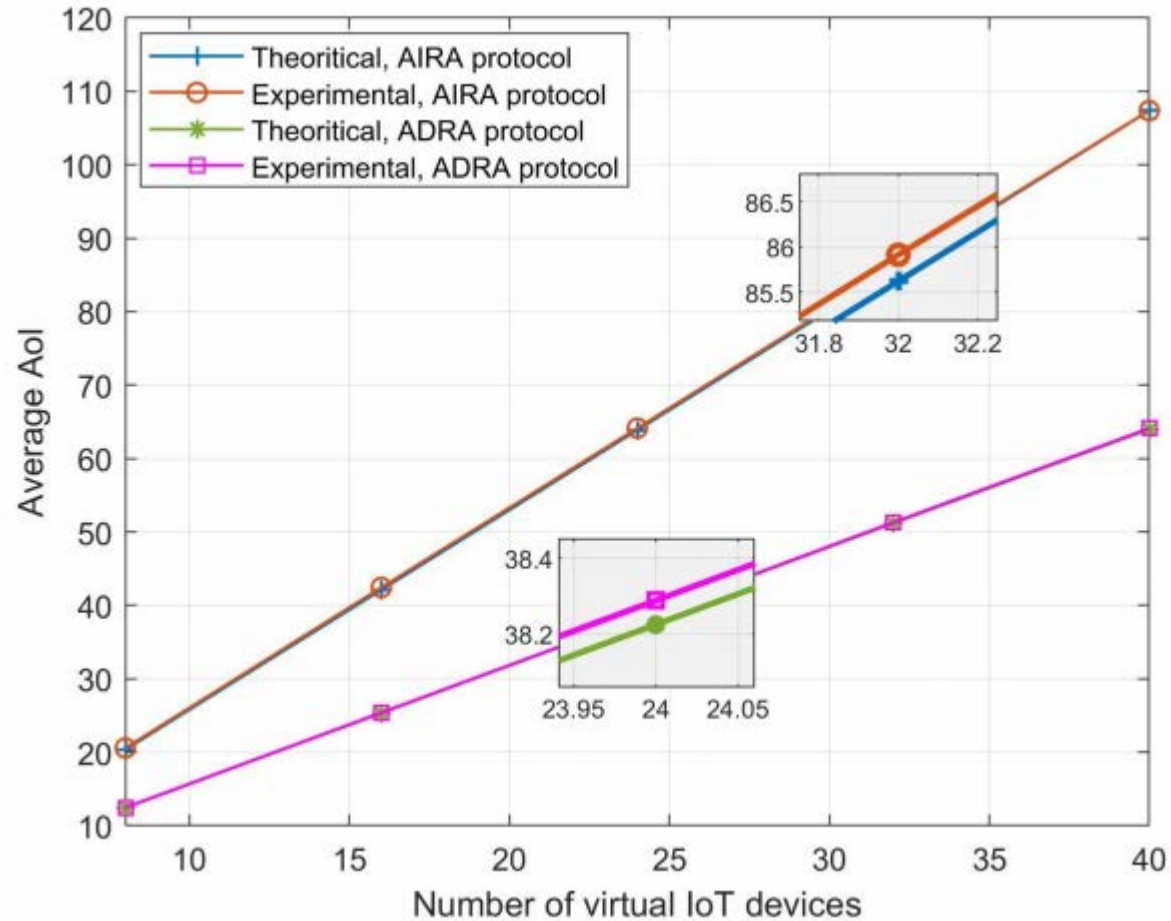
[Han'20] Z. Han, J. Liang, Y. Gu, **H. Chen**, "Software-Defined Radio Implementation of Age-of-Information-Oriented Random Access," accepted to appear in IEEE IECON, July 2020.

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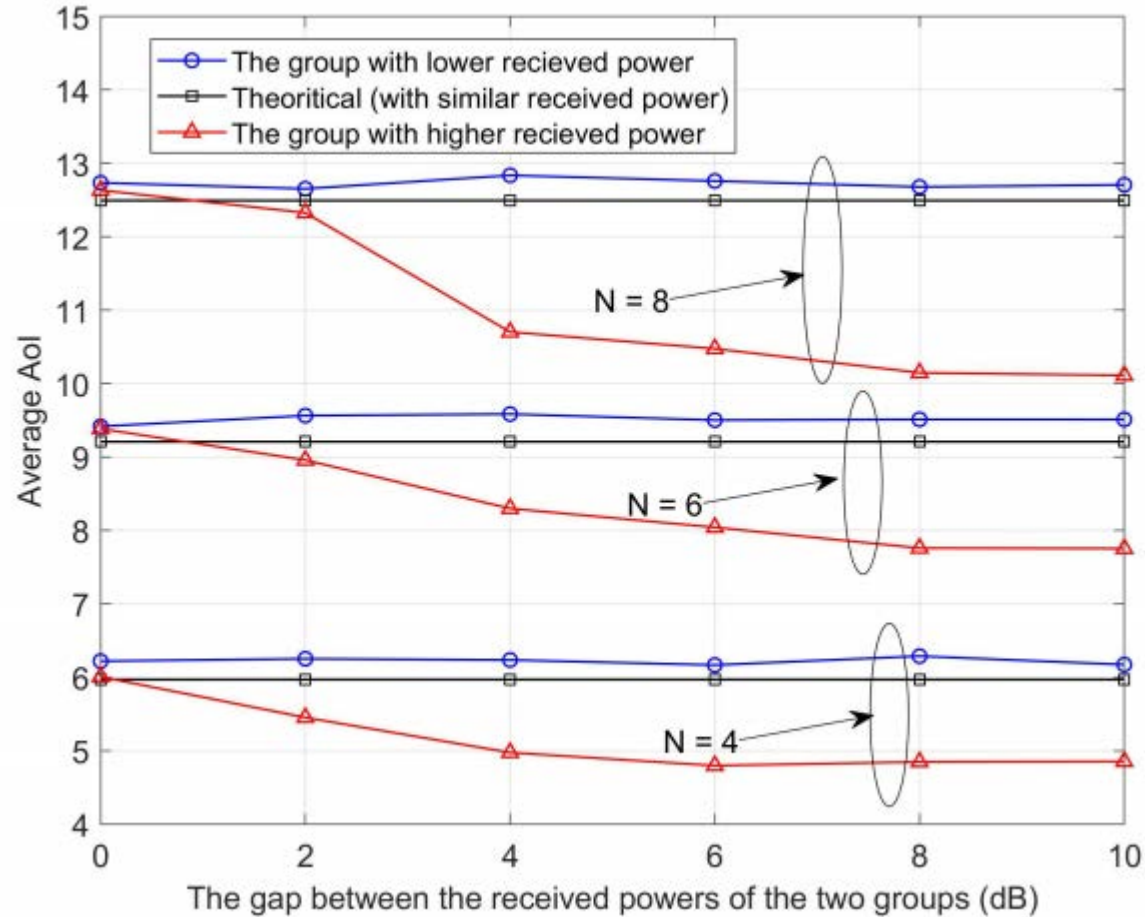
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Thanks to my students and RAs: Qian, Zixiao, Bob, Jingwei, Tong Zhang

Thanks to my collaborators for many insightful discussions: Dr Yifan Gu, Prof Yong Zhou, Prof Soung Liew, Nikos, etc.

Pass to Nikolaos PAPPAS

# Concluding remarks

Nikolaos PAPPAS

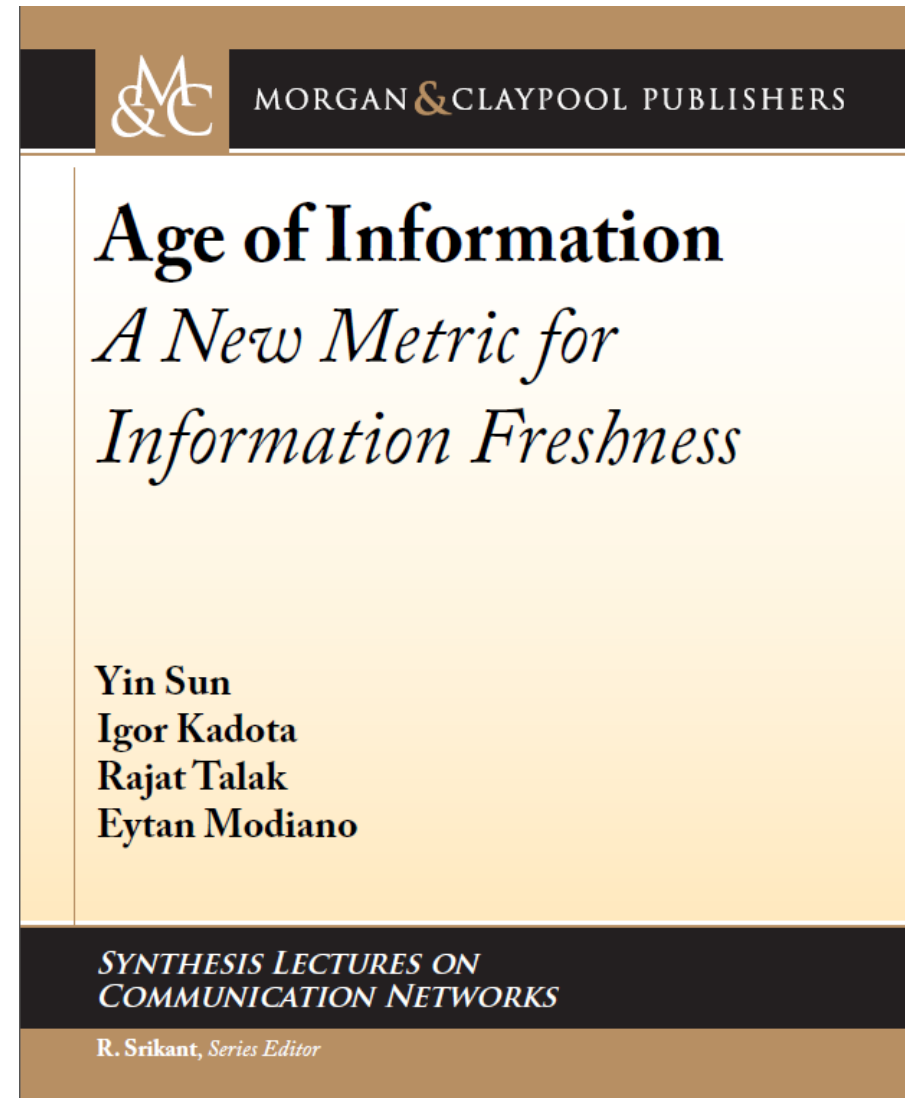
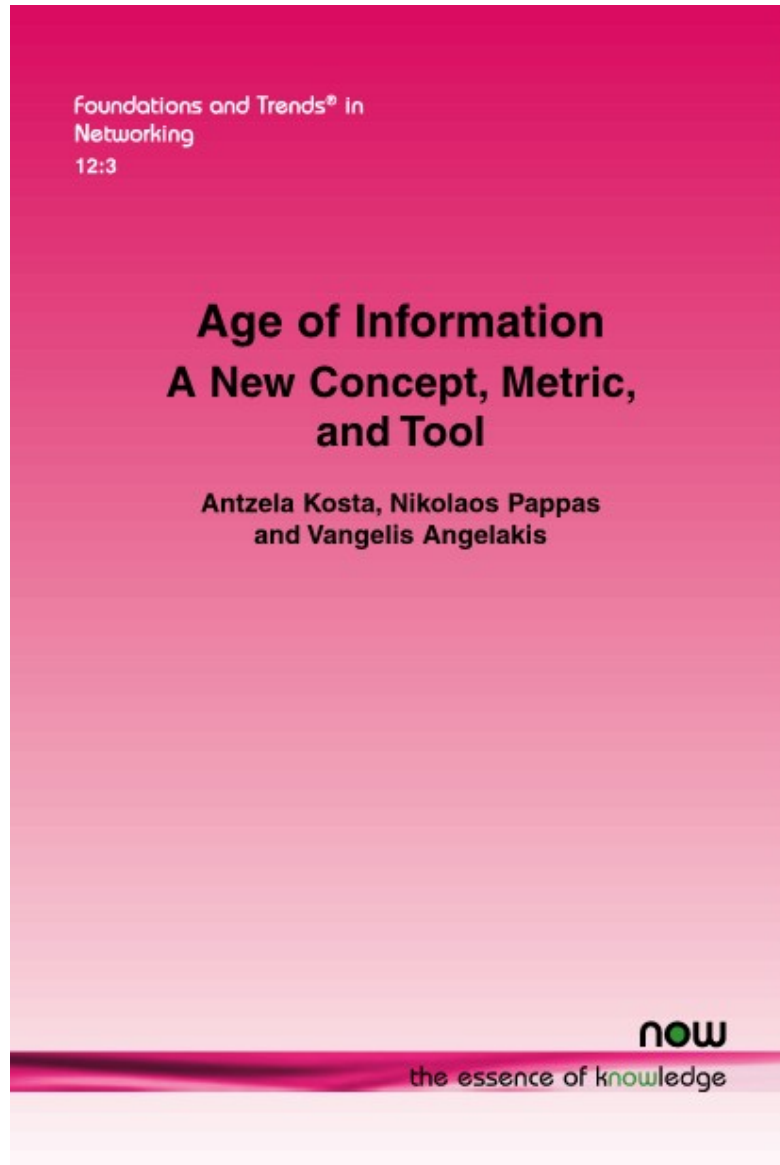
# Concluding remarks and future directions

- AoI has emerged as an end-to-end performance metric for systems that employ status updates.
- Introduction of information freshness requirements will create systems that work smarter than harder, so they will be more effective.
  - The updating process should not underload nor overload the system.
  - The system should process new updates rather than old.
  - The system should avoid processing updates without sufficient novelty.

- There are still many interesting research directions
  - Definition of effective age (term coined by Prof. Ephremides in ITA 2015)
  - Sampling and remote reconstruction
  - Deploying of AoI in machine learning
- It provides stronger connections with areas such as Signal Processing
- Metrics that can capture the requirements of Wireless Networked Control Systems
- AoI is one of the dimensions of semantics-empowered communications

M. Kountouris, N. Pappas, "Semantics-Empowered Communication for Networked Intelligent Systems", *arxiv*, July 2020.

# Monographs



# Upcoming workshop in IEEE GLOBECOM 2020



## IEEE Globecom 2020 Workshop 7: Information Freshness, Communications, Control, and Computing for Industrial IoT

Date: Monday, 7 December 2020

### Call for Workshop Papers

Industrial Internet of Things (IIoT) integrates advanced digital computing, communications, and control technology. Within the Industry 4.0 framework, IIoT aims to tackle a set of new technological challenges in industrial control, automation, and intelligence. Since most of the IIoT applications, such as intelligent transportation, telesurgery, industry automation, power systems automation, and power electronics control, are mission-critical and demand real-time communications and computation capabilities for successful closed-loop operation, IIoT requires the joint design and optimization of communications, control and computing (3C). Moreover, the freshness and value of information is significant in real-time IIoT applications.

Age of Information (AoI) has emerged as a concept that has recently gained increasing attention as a performance metric for quantifying information freshness at the destination. It is important to note that improving first-order metrics (e.g., for providing high throughput and low delay) that render enhanced performance in traditional communication networks, are inadequate when it comes to improving the information freshness. Therefore, there is an urgent need to re-examine existing metrics and develop new ones for emerging time-critical applications.

The workshop is expected to bring together academic and industrial researchers to identify and discuss the major technical challenges and recent breakthroughs related to AoI, communications, control, and computing in IIoT.

- AoI analysis and optimization
- Multiuser scheduling for optimizing information freshness
- Cooperative status update for improved information freshness
- The application of emerging technologies (e.g., cloud and edge computing, machine learning) for improving AoI performance
- AoI approaches for industrial control applications
- Fundamental tradeoffs between communications, control and computing in IIoT
- Resource allocation for next generation wireless sensing and control applications in IIoT
- MAC layer and network layer design for supporting multi-loop control applications
- Edge computing for low-latency control
- Cyber-physical security in IIoT
- Prototypes and testbeds for validating AoI and communications-control-computing codesign protocols.

### Workshop Organizers

- Yonghui Li, The University of Sydney, Australia (yonghui.li@sydney.edu.au)
- He (Henry) Chen, The Chinese University of Hong Kong, Hong Kong SAR, China (he.chen@ie.cuhk.edu.hk)
- Nikolaos Pappas, Linköping University, Sweden (nikolaos.pappas@liu.se)
- Sheng Zhou, Tsinghua University, China (sheng.zhou@tsinghua.edu.cn)
- Zhibo Pang, ABB AB, Corporate Research, Sweden (pang.zhibo@se.abb.com)

### IMPORTANT DATES

Paper Submission Deadline: ~~1 August 2020~~ 14 August 2020 (Firm)

Paper Acceptance Notification: 15 September 2020

Camera-Ready: 1 October 2020

Thanks for your attention.

Q&A

[www.liu.se](http://www.liu.se)